First-Order Logic

Recap from Last Time

Recap So Far

- A *propositional variable* is a variable that is either true or false.
- The *propositional connectives* are as follows:
 - Negation: $\neg p$
 - Conjunction: $p \land q$
 - Disjunction: *p* V *q*
 - Implication: $p \rightarrow q$
 - Biconditional: $p \leftrightarrow q$
 - True: ⊤
 - False: \bot

Take out a sheet of paper!

What's the truth table for the \rightarrow connective?

What's the negation of $p \rightarrow q$?

New Stuff!

First-Order Logic

What is First-Order Logic?

- **First-order logic** is a logical system for reasoning about properties of objects.
- Augments the logical connectives from propositional logic with
 - predicates that describe properties of objects,
 - *functions* that map objects to one another, and
 - *quantifiers* that allow us to reason about multiple objects.

Some Examples

$Likes(You, Eggs) \land Likes(You, Tomato) \rightarrow Likes(You, Shakshuka)$



> These blue terms are called *constant symbols.* Unlike propositional variables, they refer to *objects*, not *propositions*.

> The red things that look like function calls are called *predicates*. Predicates take objects as arguments and evaluate to true or false.

> What remains are traditional propositional connectives. Because each predicate evaluates to true or false, we can connect the truth values of predicates using normal propositional connectives.

Reasoning about Objects

- To reason about objects, first-order logic uses predicates.
- Examples:

Cute(Quokka)

Cool(CS103 students)

- Applying a predicate to arguments produces a proposition, which is either true or false.
- Typically, when you're working in FOL, you'll have a list of predicates, what they stand for, and how many arguments they take. It'll be given separately from the formulas you write.

First-Order Sentences

 Sentences in first-order logic can be constructed from predicates applied to objects:
 Cute(a) → Dikdik(a) ∨ Kitty(a) ∨ Puppy(a)

 $Succeeds(You) \leftrightarrow Practices(You)$

 $x < 8 \rightarrow x < 137$

The less-than sign is just another predicate. Binary predicates are sometimes written in *infix notation* this way.

Numbers are not "built in" to first-order logic. They're constant symbols just like "You" and "a" above.

Equality

- First-order logic is equipped with a special predicate = that says whether two objects are equal to one another.
- Equality is a part of first-order logic, just as \rightarrow and \neg are.
- Examples:

MilesMorales = SpiderMan MorningStar = EveningStar

• Equality can only be applied to **objects**; to state that two **propositions** are equal, use \leftrightarrow .

Let's see some more examples.

These purple terms are *functions*. Functions take objects as input and produce objects as output.

Functions

- First-order logic allows *functions* that return objects associated with other objects.
- Examples:

ColorOf(Money) MedianOf(x, y, z) x + y

- As with predicates, functions can take in any number of arguments, but always return a single value.
- Functions evaluate to *objects*, not *propositions*.

Objects and Predicates

- When working in first-order logic, be careful to keep objects (actual things) and propositions (true or false) separate.
- You cannot apply functions to propositions:

• Ever get confused? *Just ask!*

The Type-Checking Table

	operate on	and produce
Connectives $(\leftrightarrow, \Lambda, \text{ etc.}) \dots$	propositions	a proposition
Predicates (=, etc.)	objects	a proposition
Functions	objects	an object

One last (and major) change

Some spider is radioactive.

Some spider is radioactive. ∃s. (Spider(s) ∧ Radioactive(s))

Some spider is radioactive.

∃s. (Spider(s) ∧ Radioactive(s))
∃ is the existential quantifier
and says "for some choice of s, the following is true."

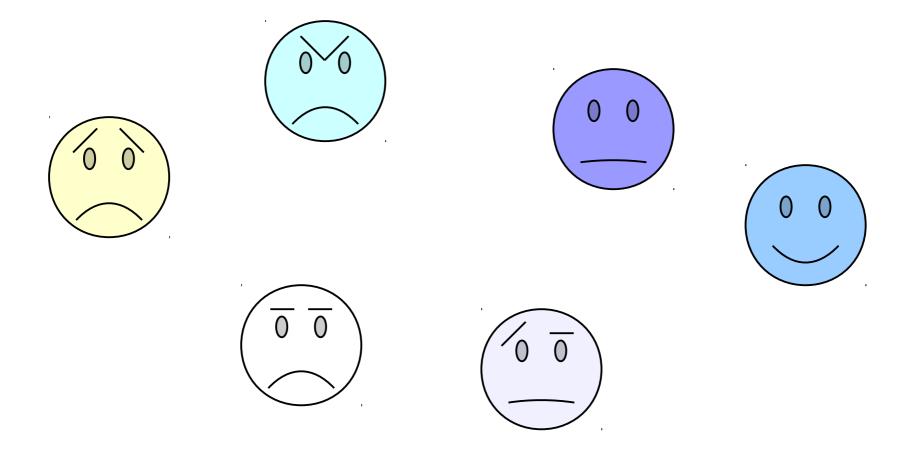
• A statement of the form

∃x. some-formula

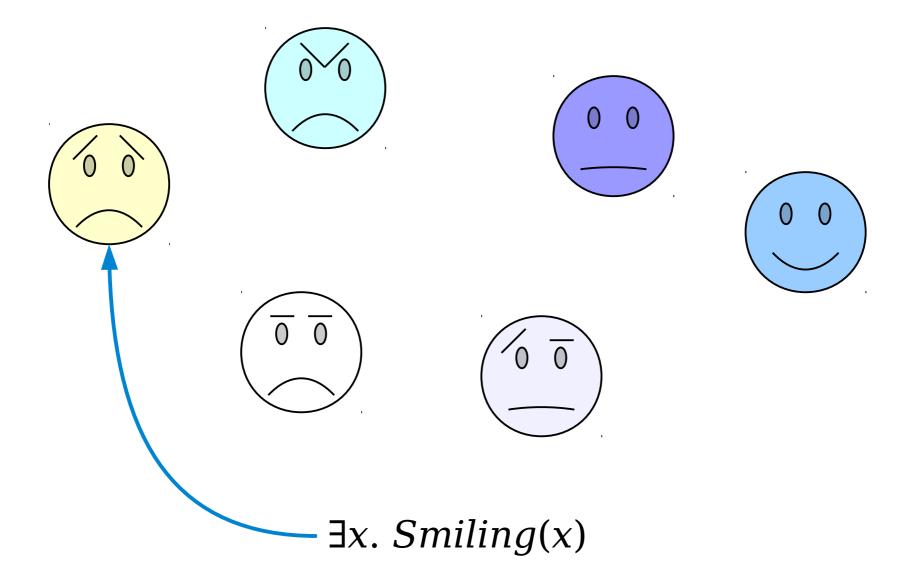
is true if there exists a choice of *x* where **some-formula** is true when that *x* is plugged into it.

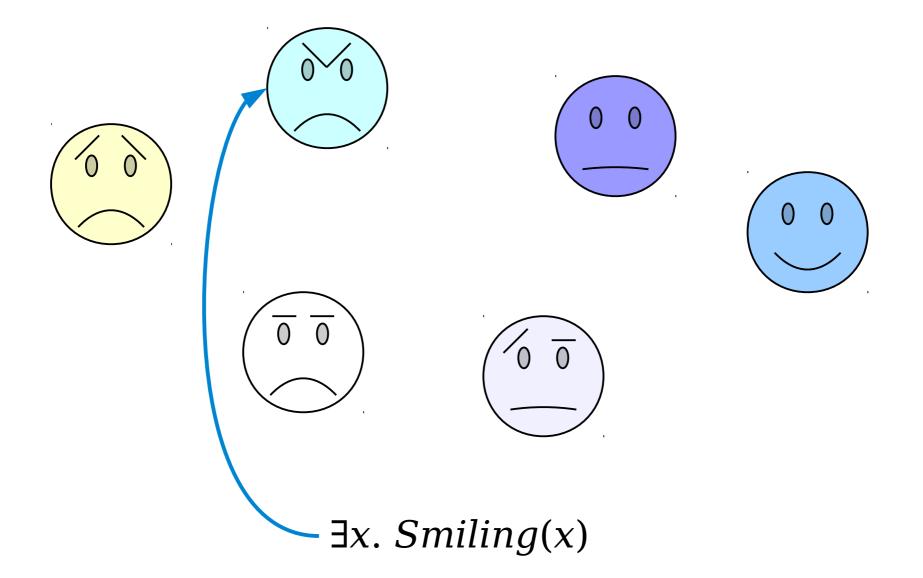
• Examples:

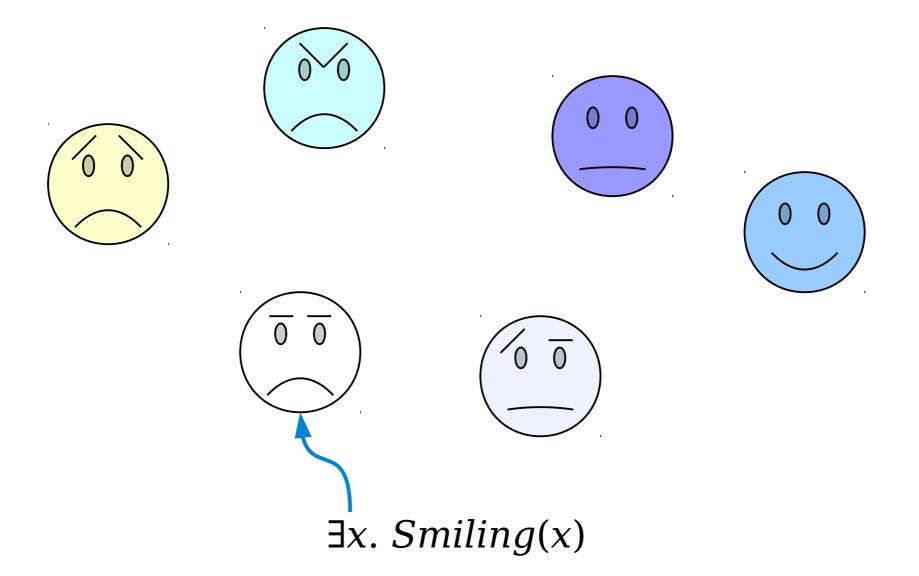
 $\exists x. (Even(x) \land Prime(x))$ $\exists x. (TallerThan(x, me) \land LighterThan(x, me))$ $(\exists w. Will(w)) → (\exists x. Way(x))$

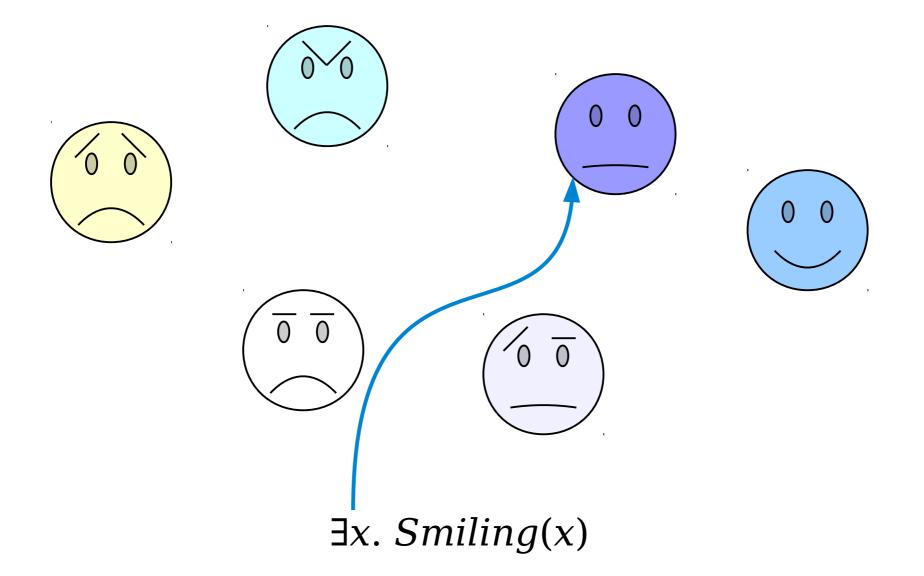


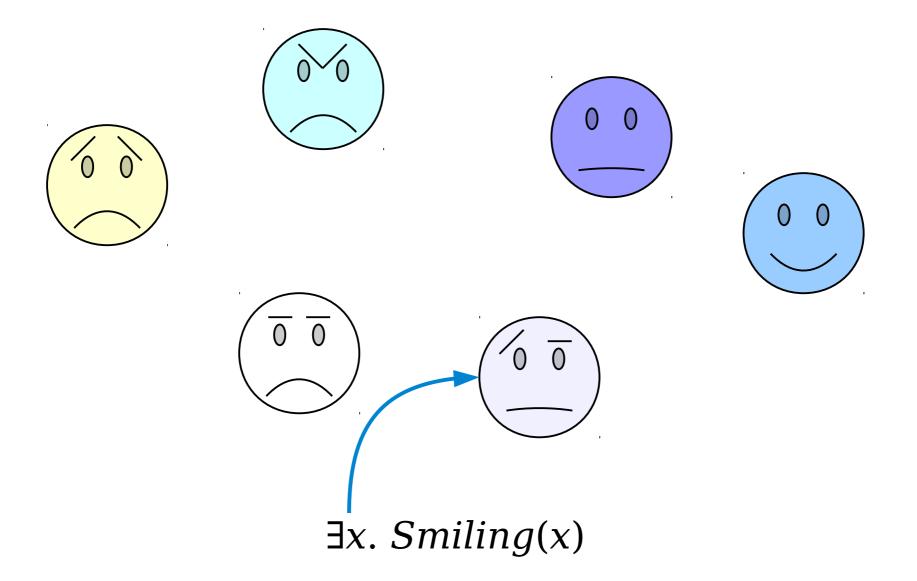
 $\exists x. Smiling(x)$

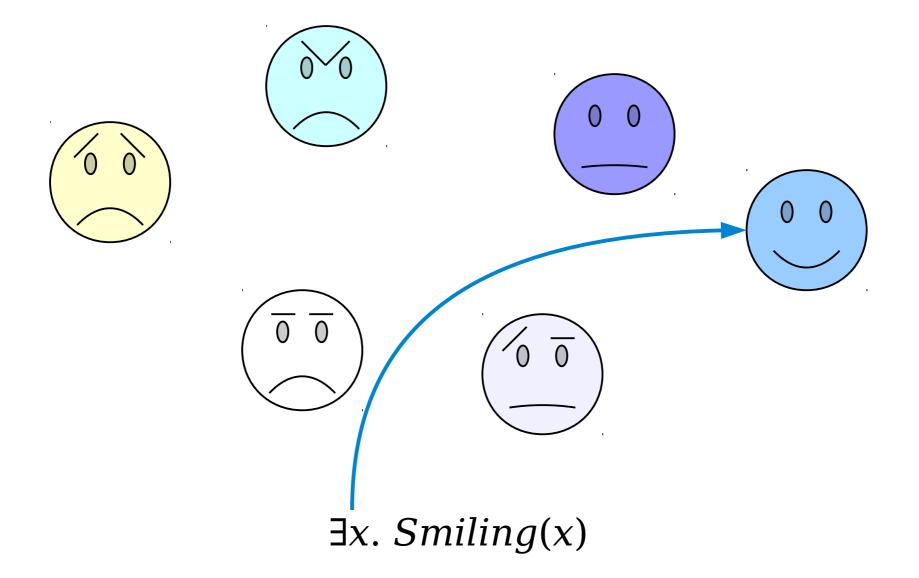


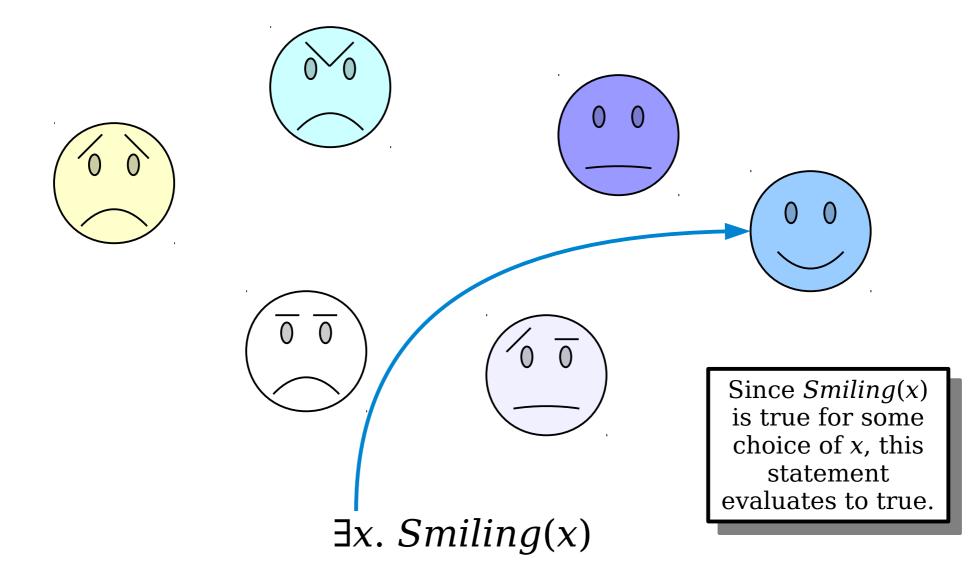


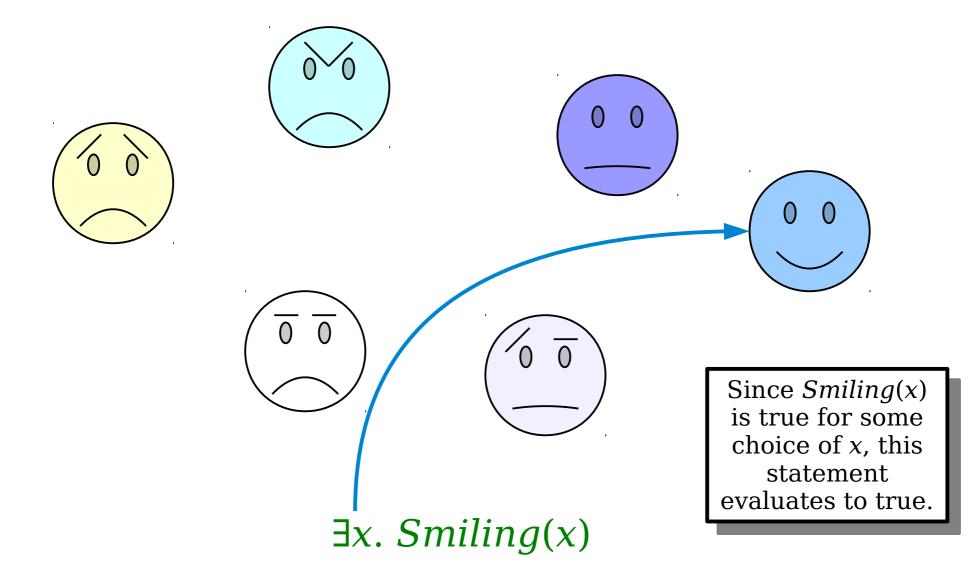


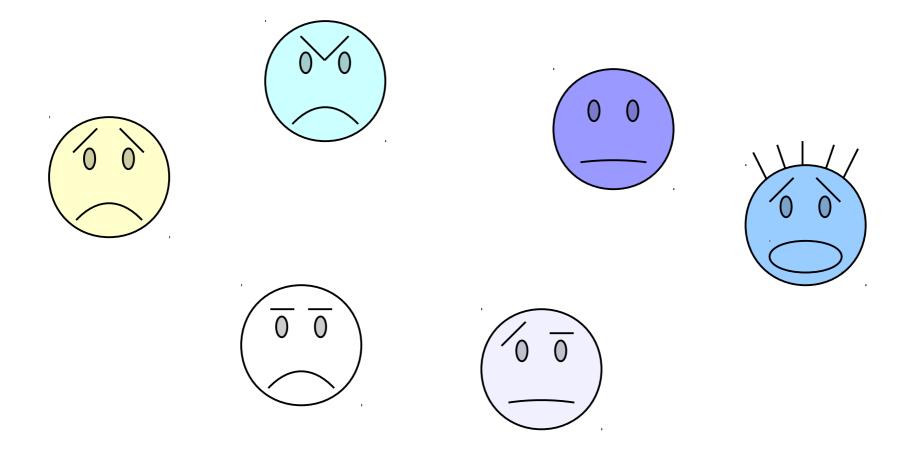




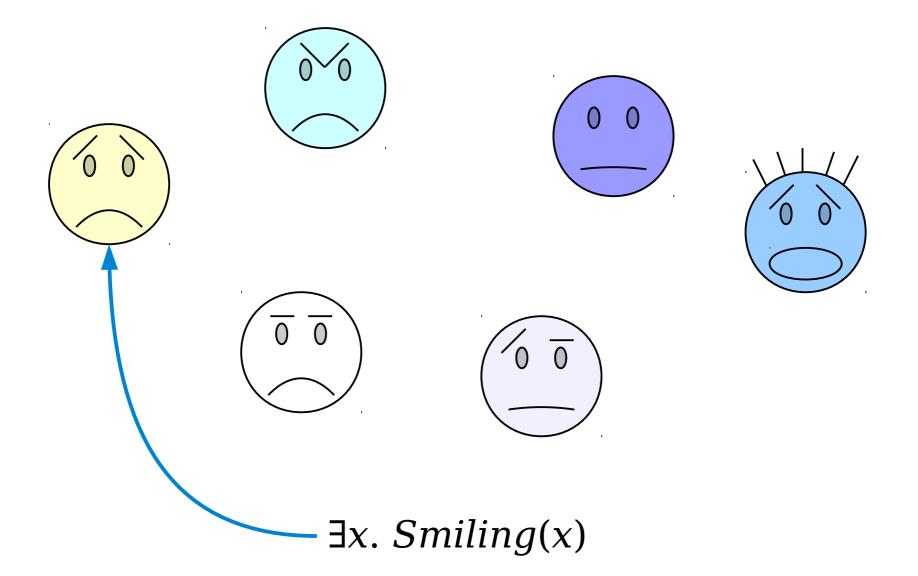


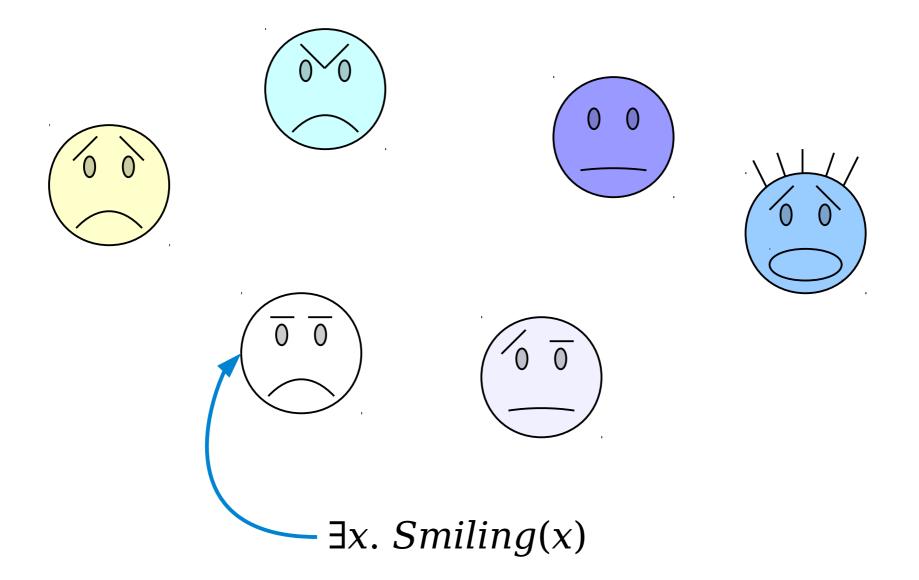


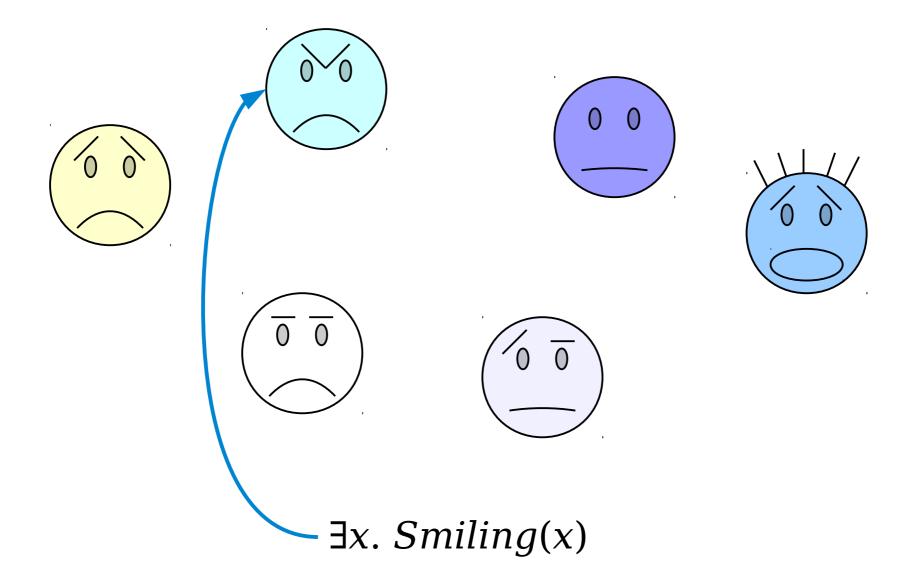


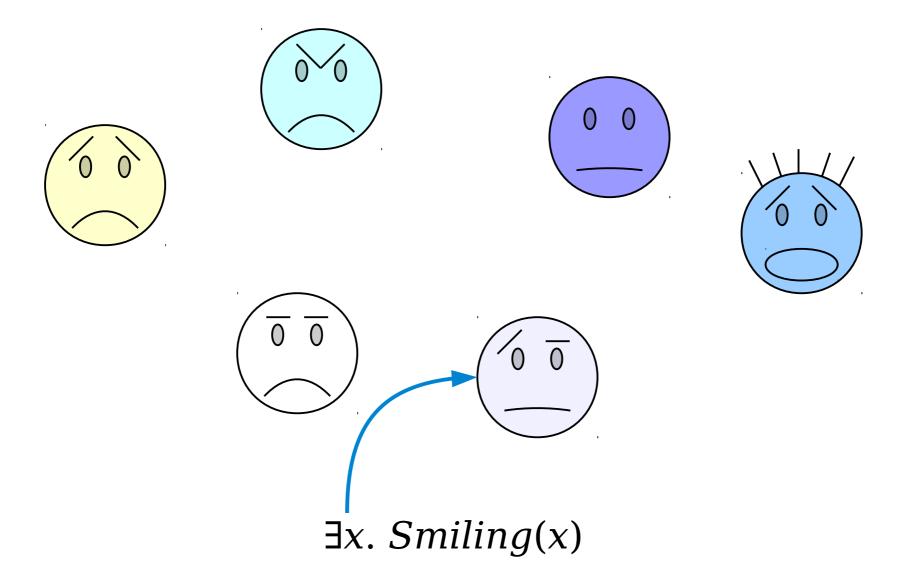


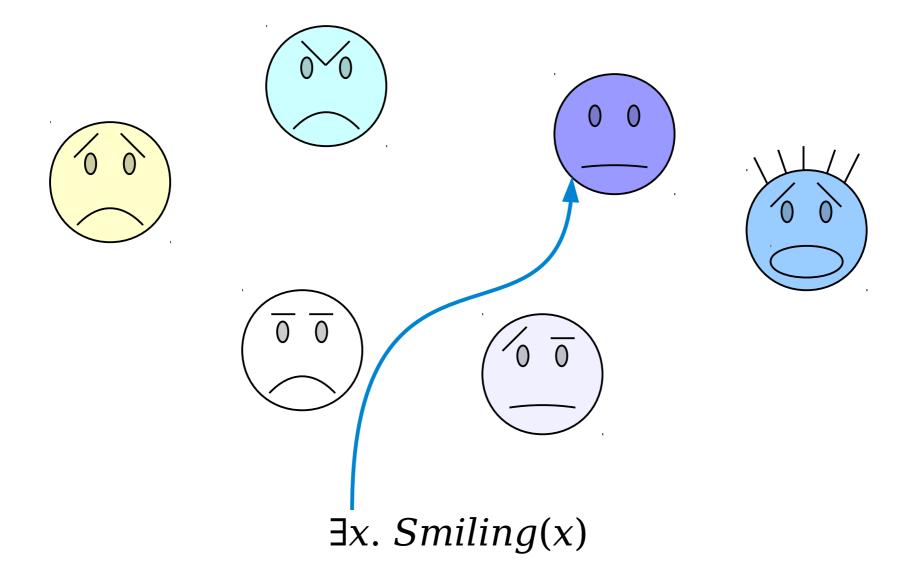
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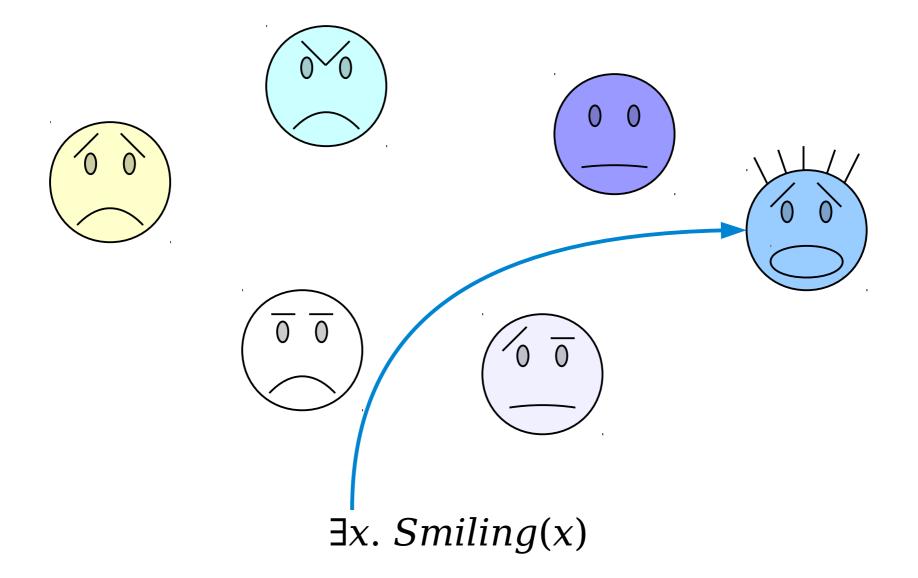


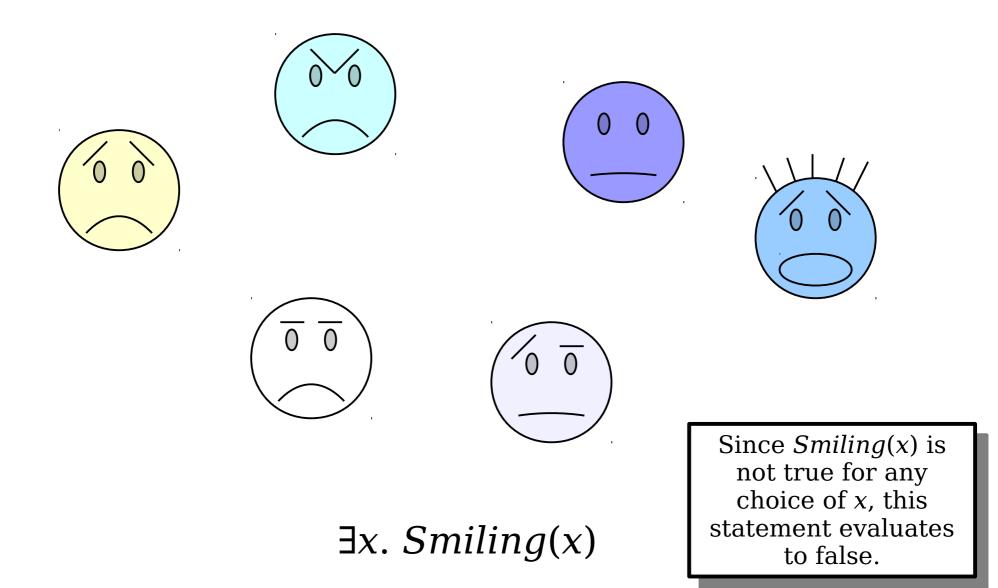


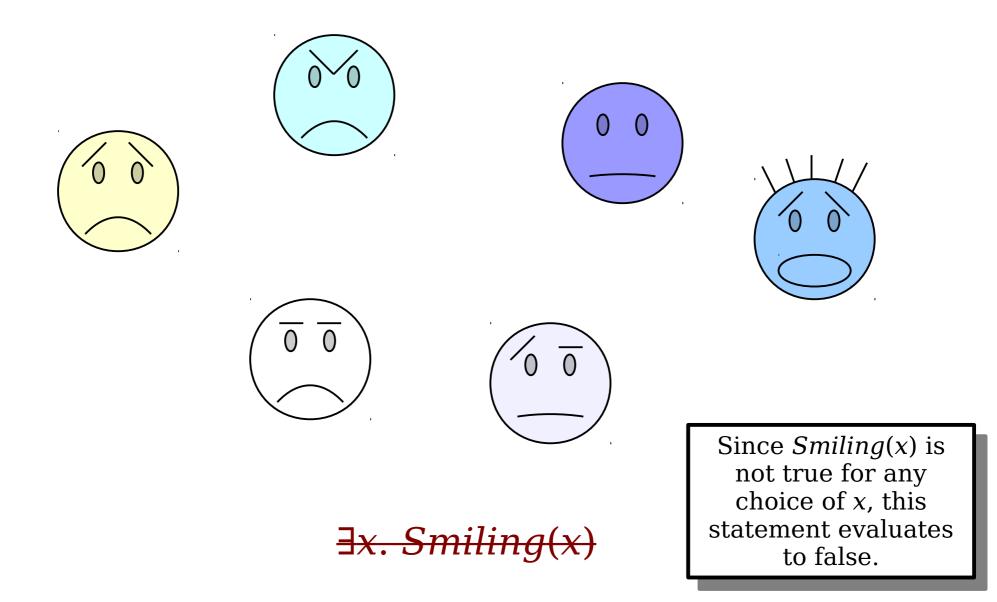






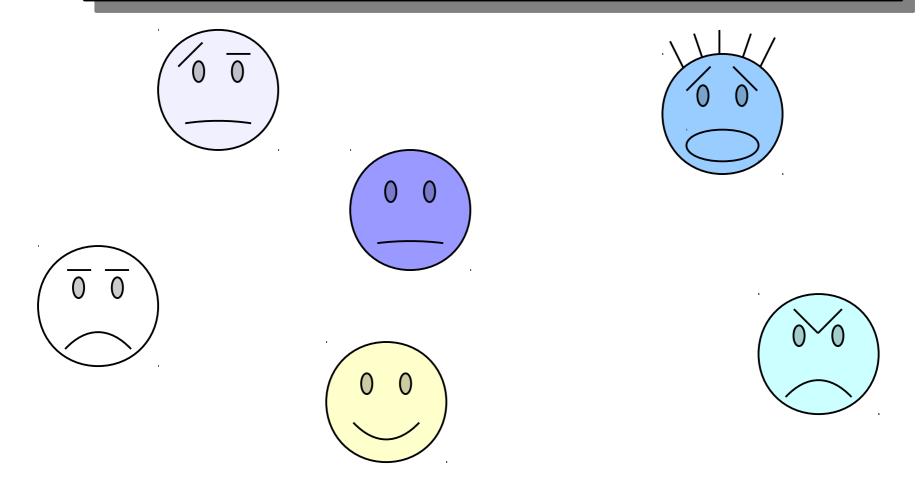


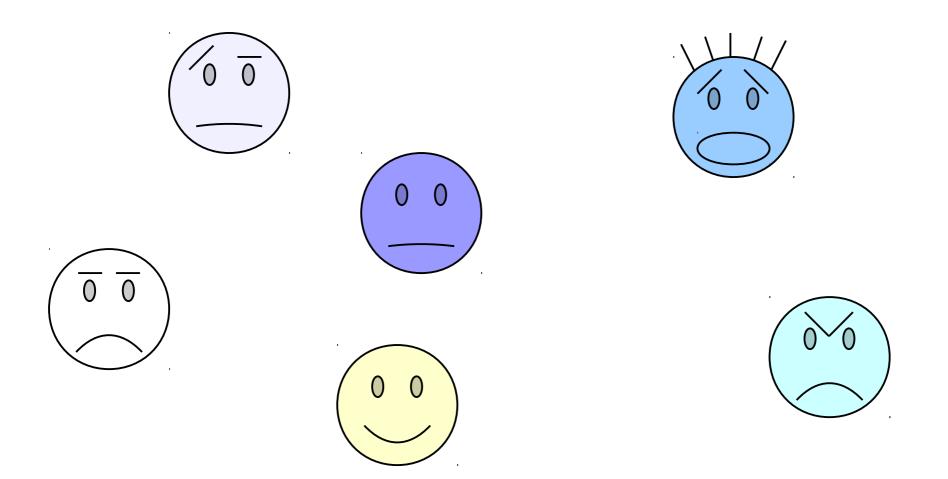


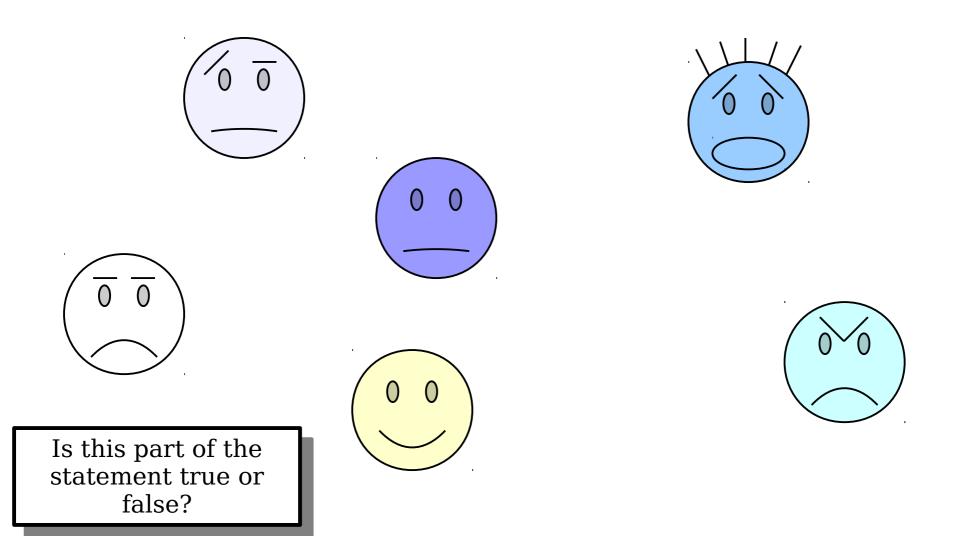


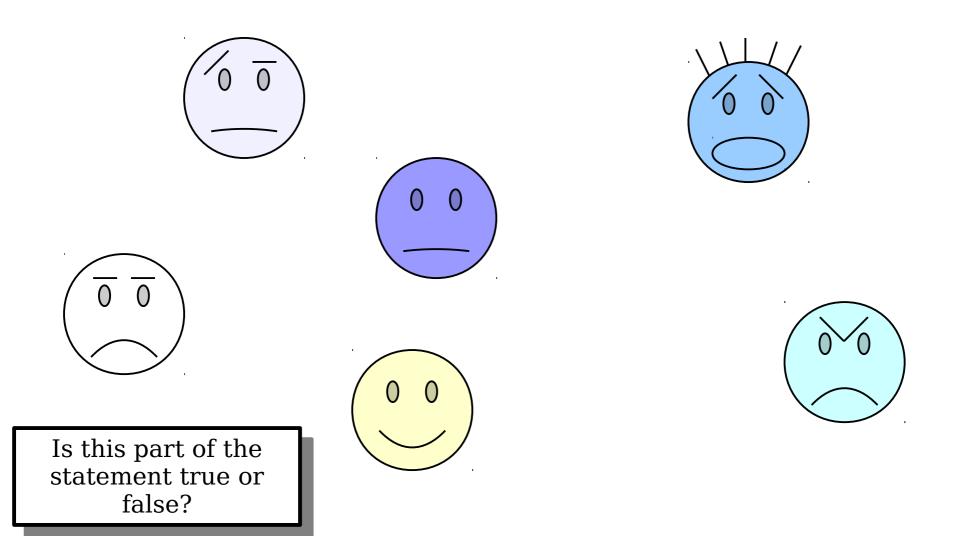
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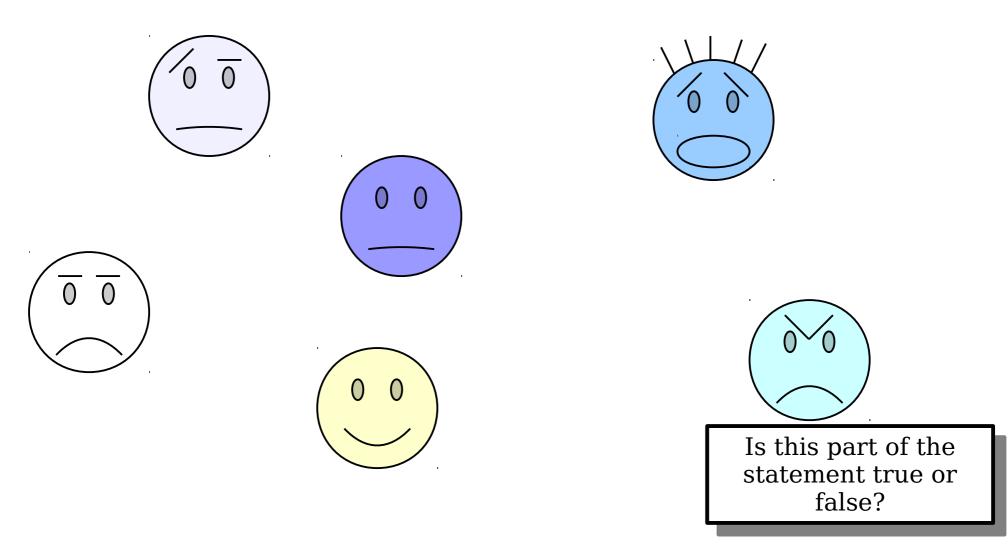
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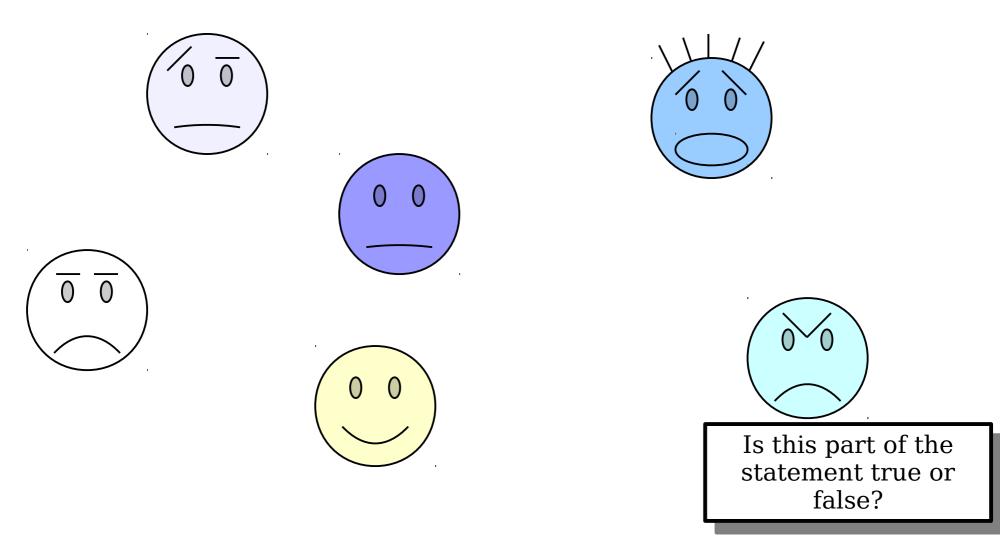


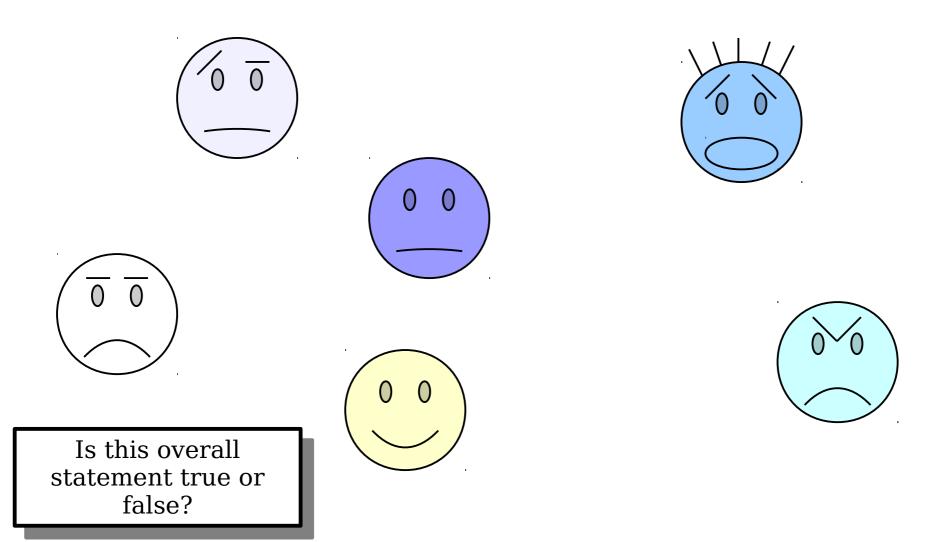


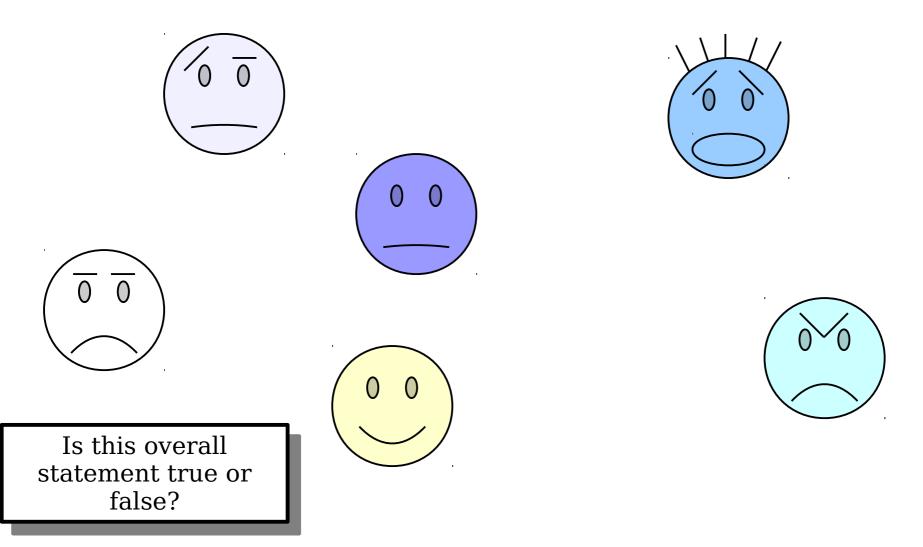












Fun with Edge Cases

 $\exists x. Smiling(x)$

Fun with Edge Cases

Existentially-quantified statements are false in an empty world, since nothing exists, period!

∃x. *Smiling*(x)

Some Technical Details

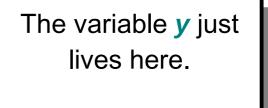
- Each quantifier has two parts:
 - the variable that is introduced, and
 - the statement that's being quantified.
- The variable introduced is scoped just to the statement being quantified.

 $(\exists x. Loves(You, x)) \land (\exists y. Loves(y, You))$

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 $(\exists x. Loves(You, x)) \land (\exists x. Loves(x, You))$

The variable **x** just lives here.

A different variable, also named **x**, just lives here.

Operator Precedence (Again)

- When writing out a formula in first-order logic, quantifiers have precedence just below ¬.
- The statement

 $\exists x. \ P(x) \land R(x) \land Q(x)$

is parsed like this:

 $\triangle \qquad (\exists x. P(x)) \land (R(x) \land Q(x)) \qquad \triangle$

- This is syntactically invalid because the variable x is out of scope in the back half of the formula.
- To ensure that x is properly quantified, explicitly put parentheses around the region you want to quantify:

 $\exists x. (P(x) \land R(x) \land Q(x))$

"For any natural number n, n is even if and only if n^2 is even"

"For any natural number *n*, *n* is even if and only if *n*² is even"

 $\forall n. (n \in \mathbb{N} \rightarrow (Even(n) \leftrightarrow Even(n^2)))$

"For any natural number n, n is even if and only if n^2 is even"

 $\forall n. \ (n \in \mathbb{N} \rightarrow (Even(n) \leftrightarrow Even(n^2)))$

∀ is the *universal quantifier* and says "for any choice of *n*, the following is true."

• A statement of the form

$\forall x. some-formula$

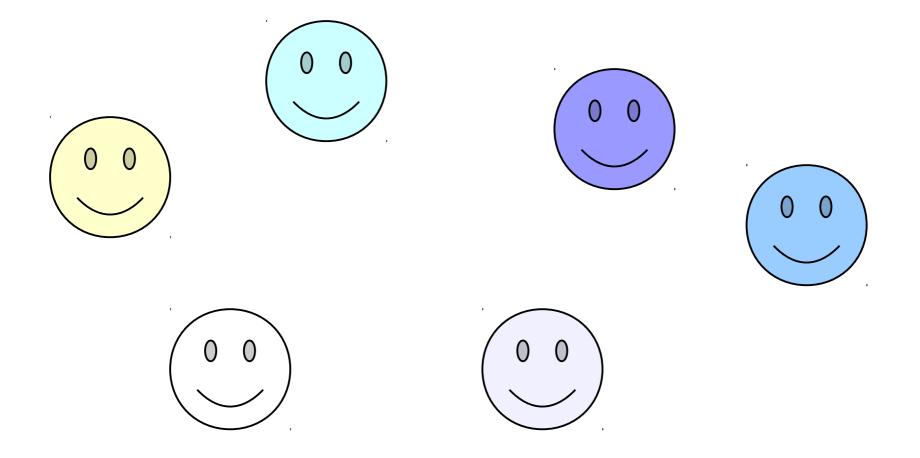
is true if, for every choice of *x*, the statement **some-formula** is true when *x* is plugged into it.

- Examples:
 - $\forall p. \ (Puppy(p) \rightarrow Cute(p))$

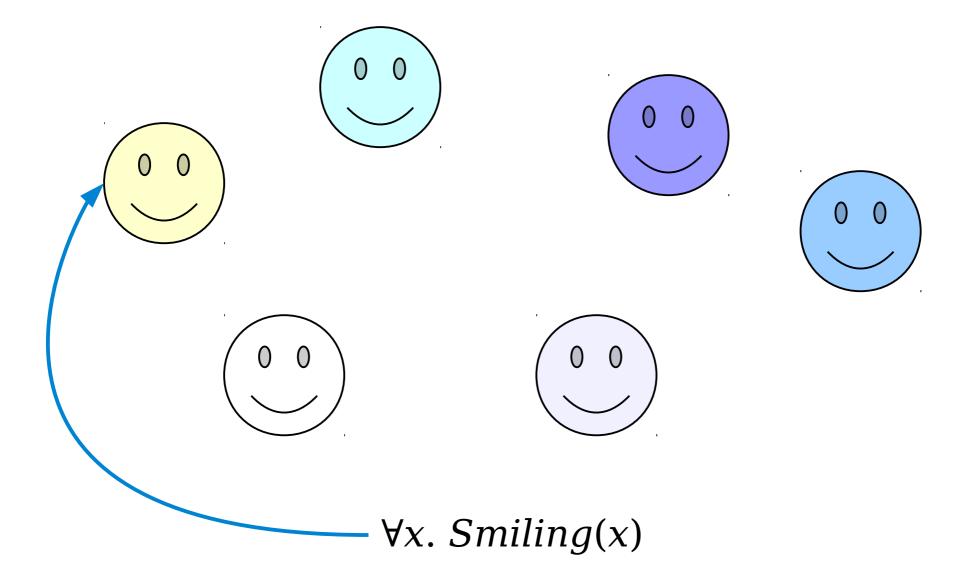
 $\forall a. (EatsPlants(a) \lor EatsAnimals(a))$

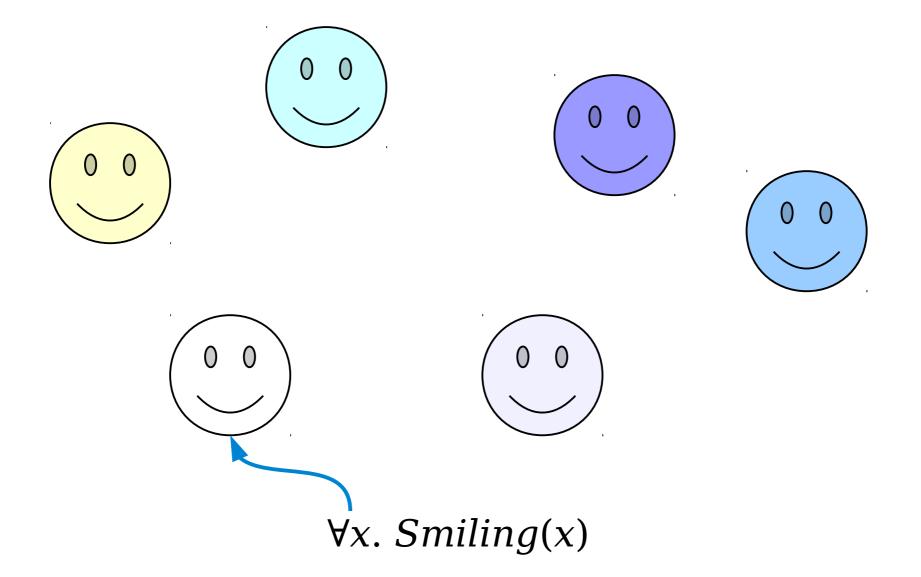
 $Tallest(SultanK\"osen) \rightarrow$

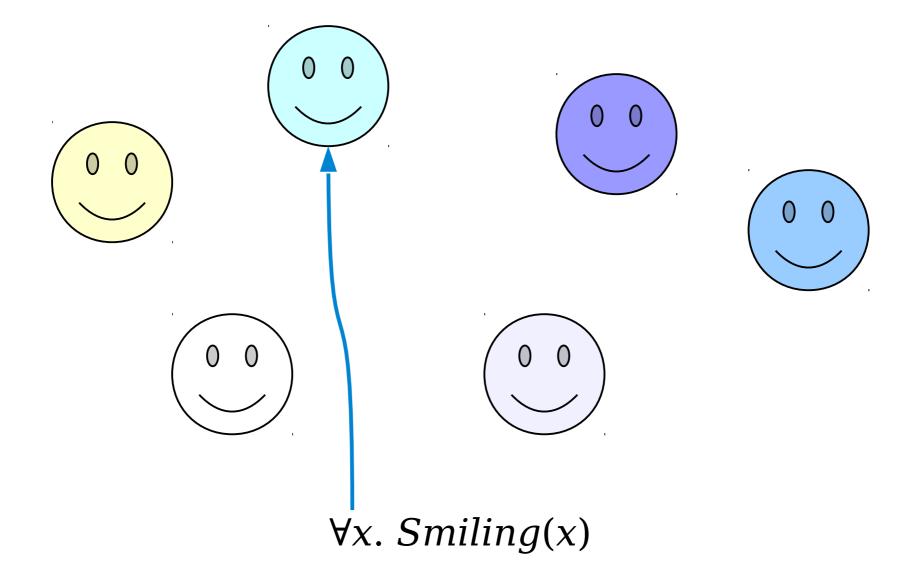
 $\forall x. (SultanK \ddot{o}sen \neq x \rightarrow ShorterThan(x, SultanK \ddot{o}sen))$

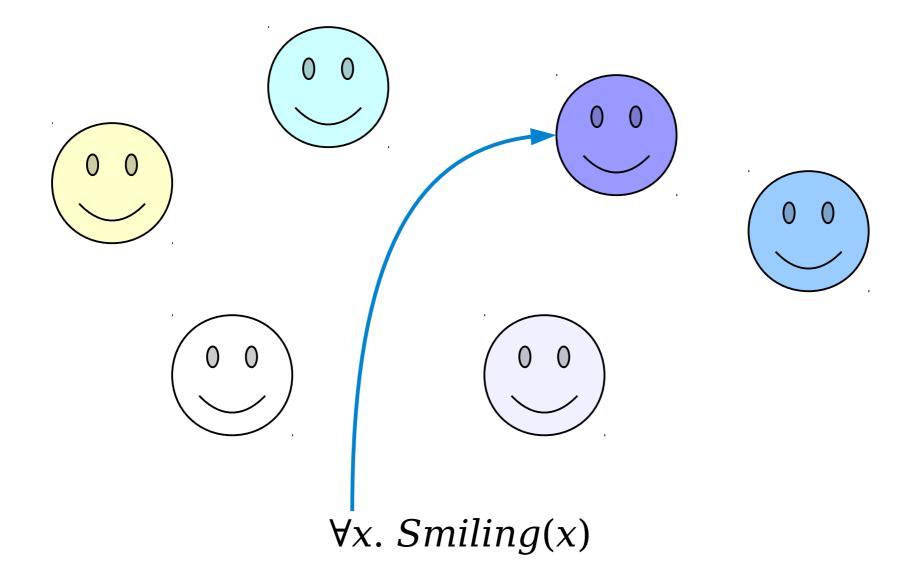


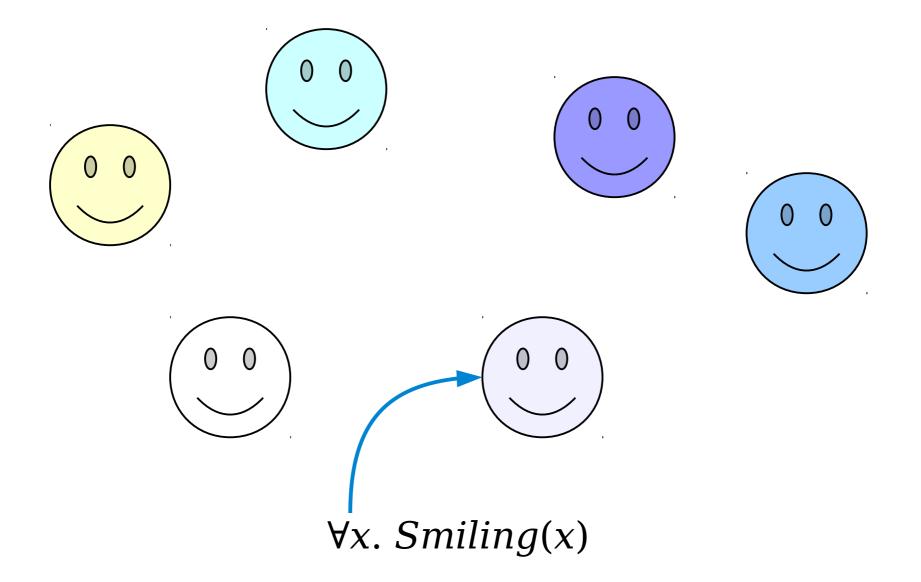
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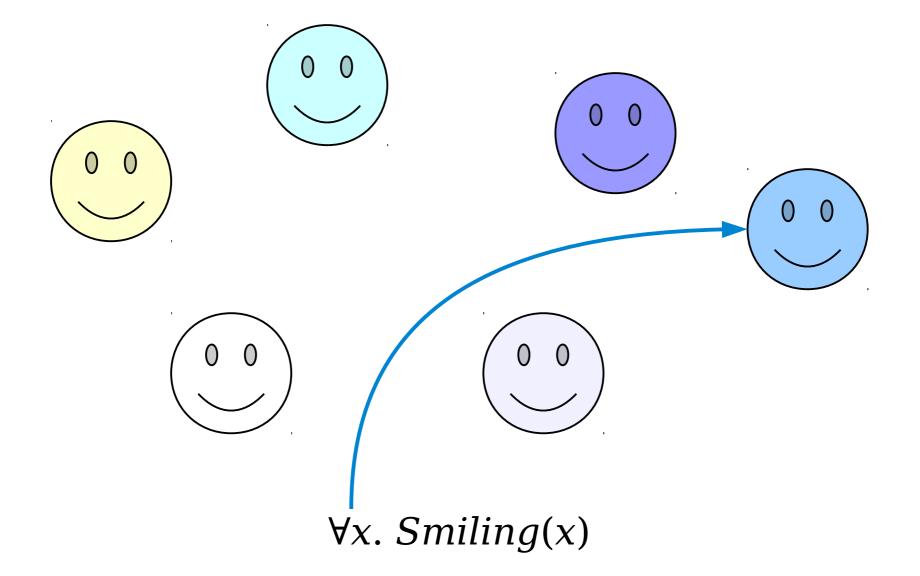


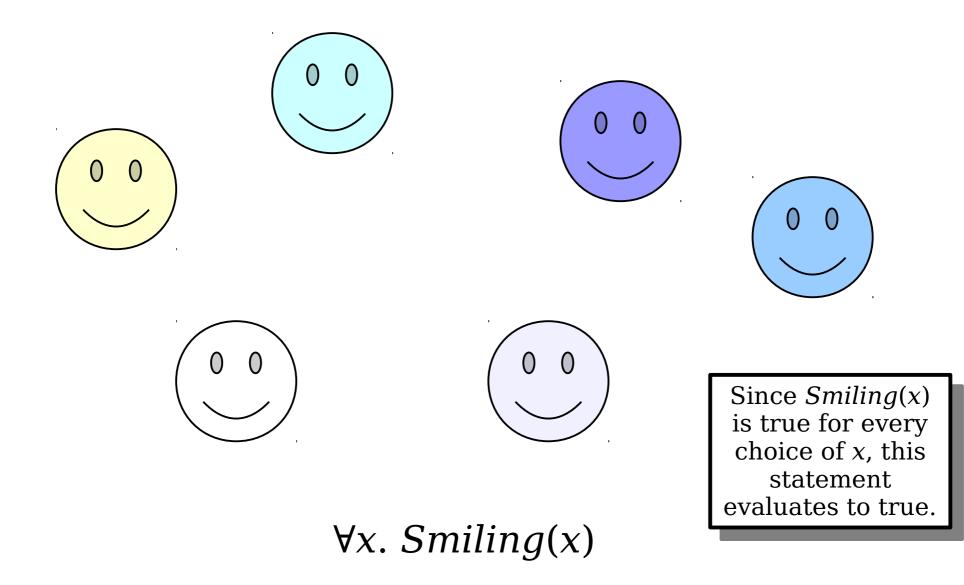


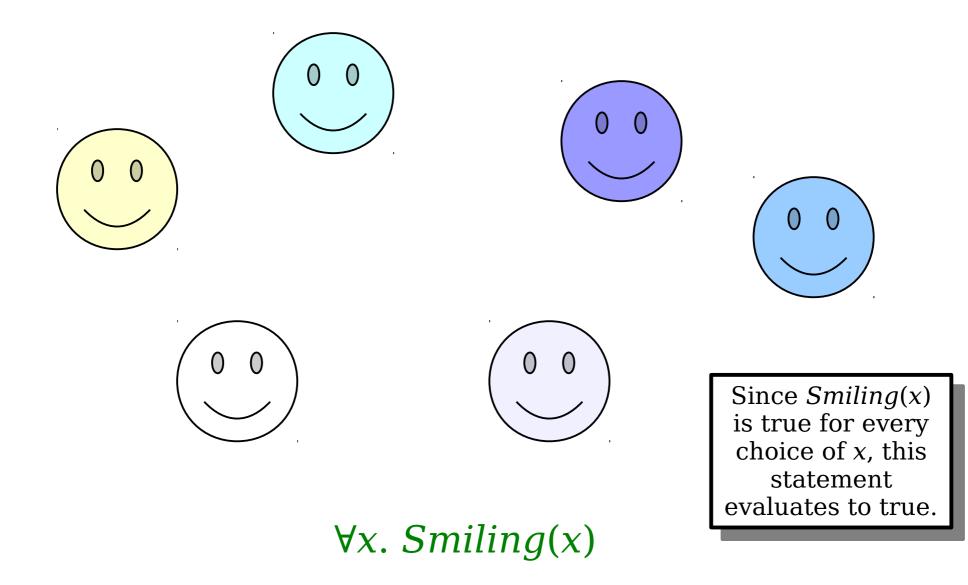


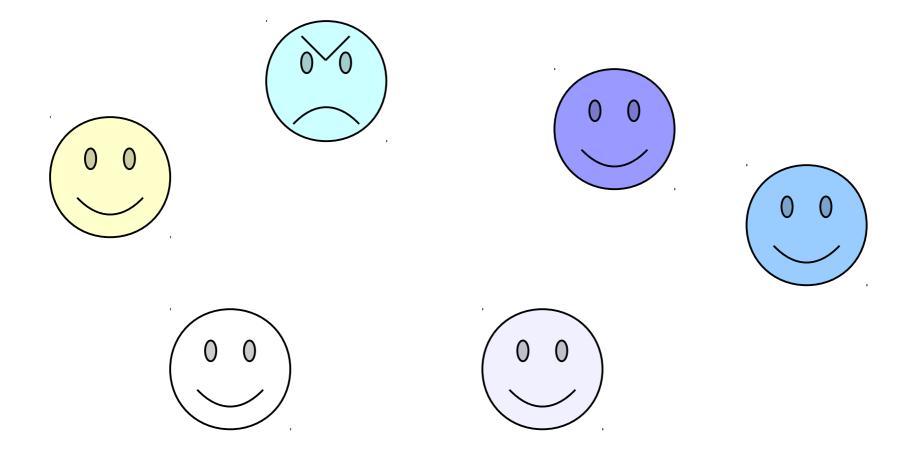




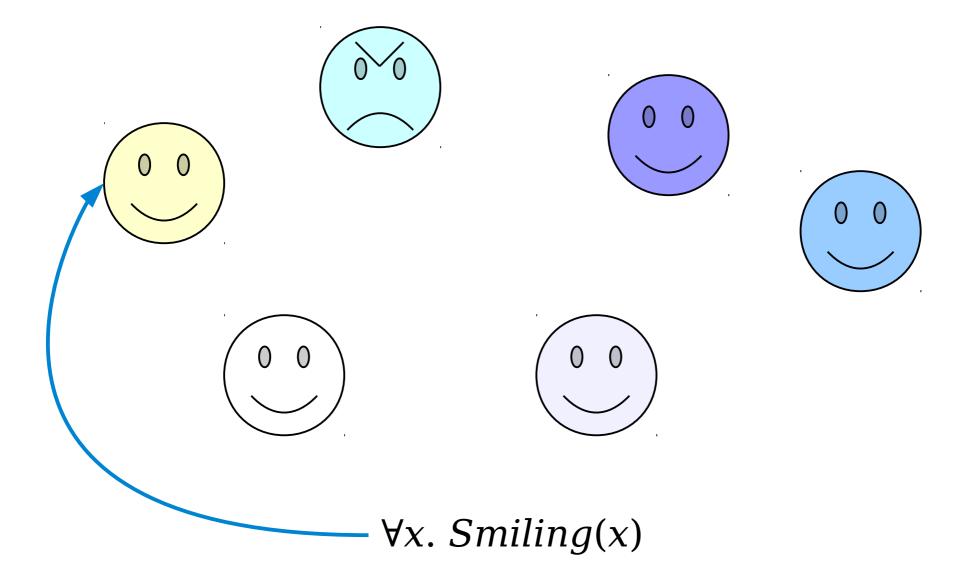


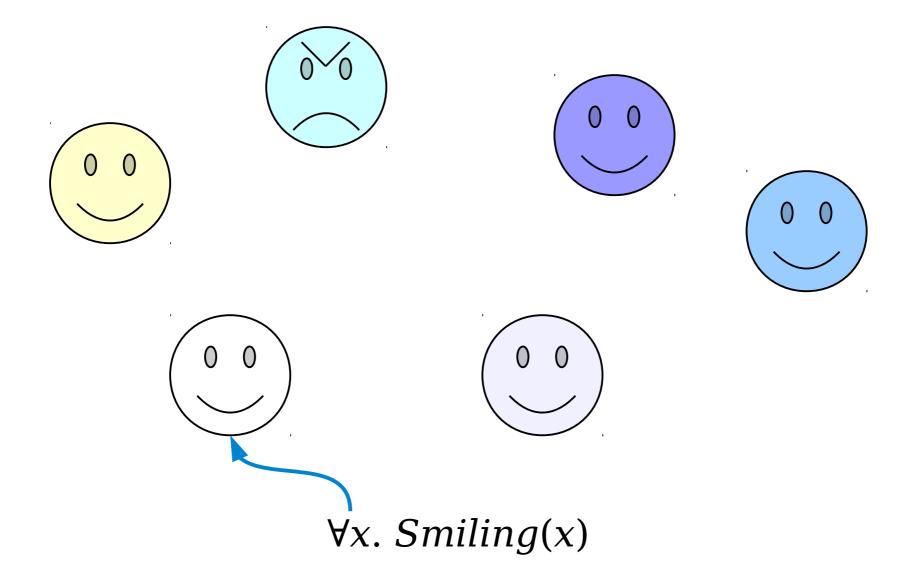


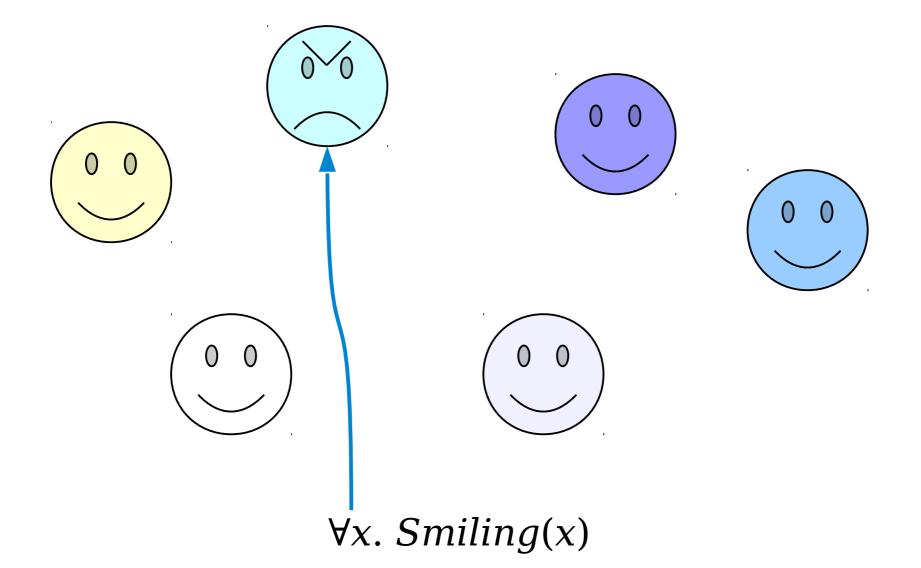


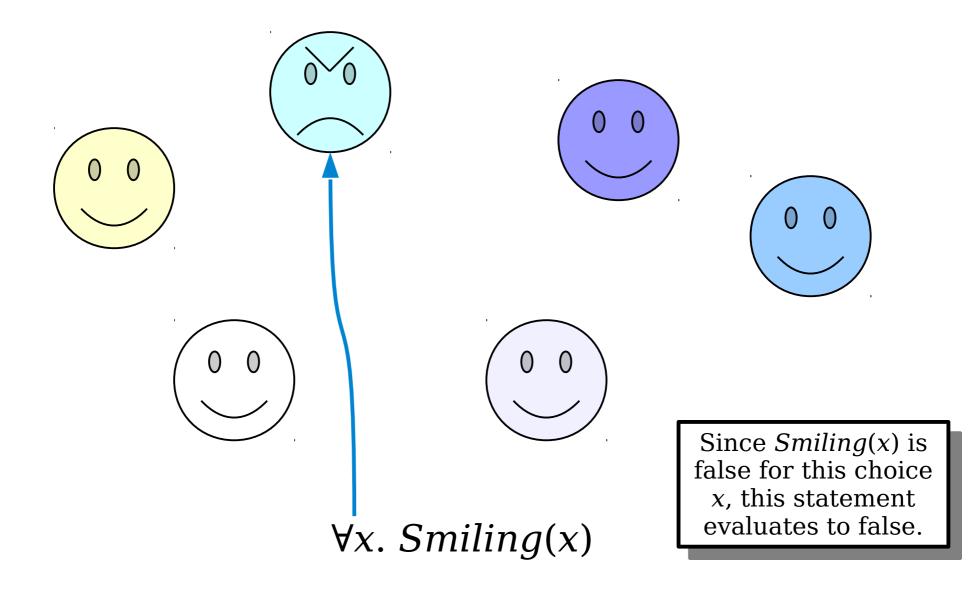


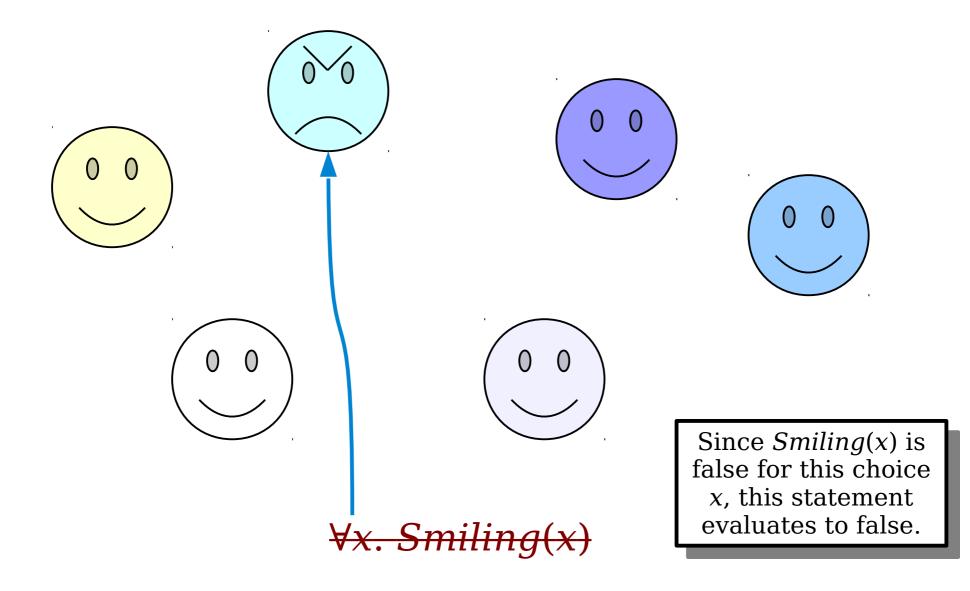
 $\forall x. Smiling(x)$





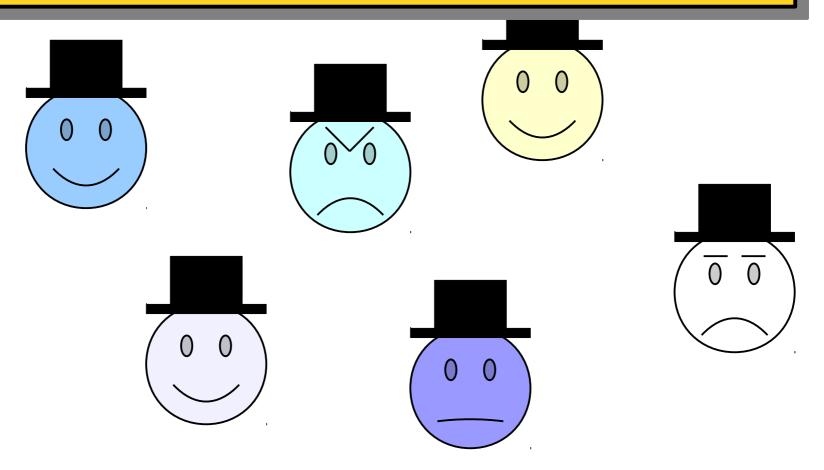


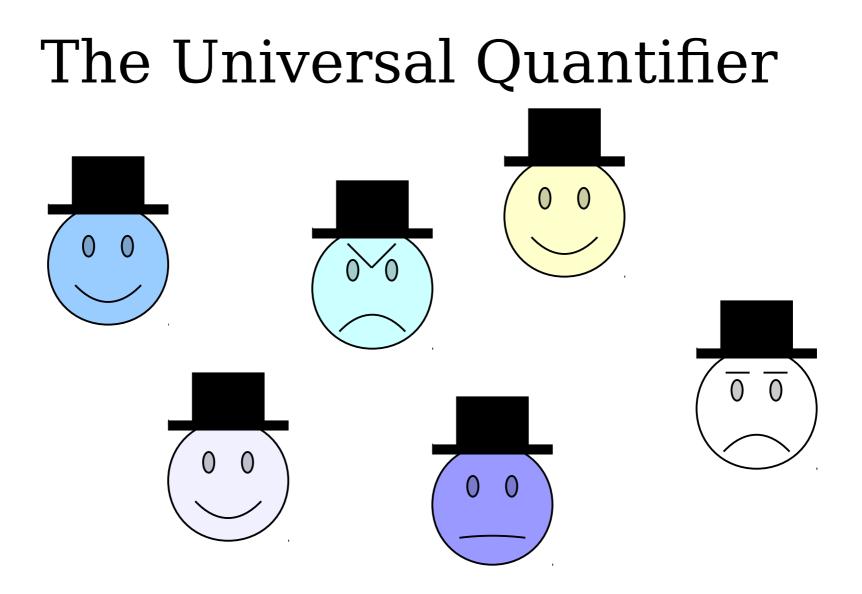


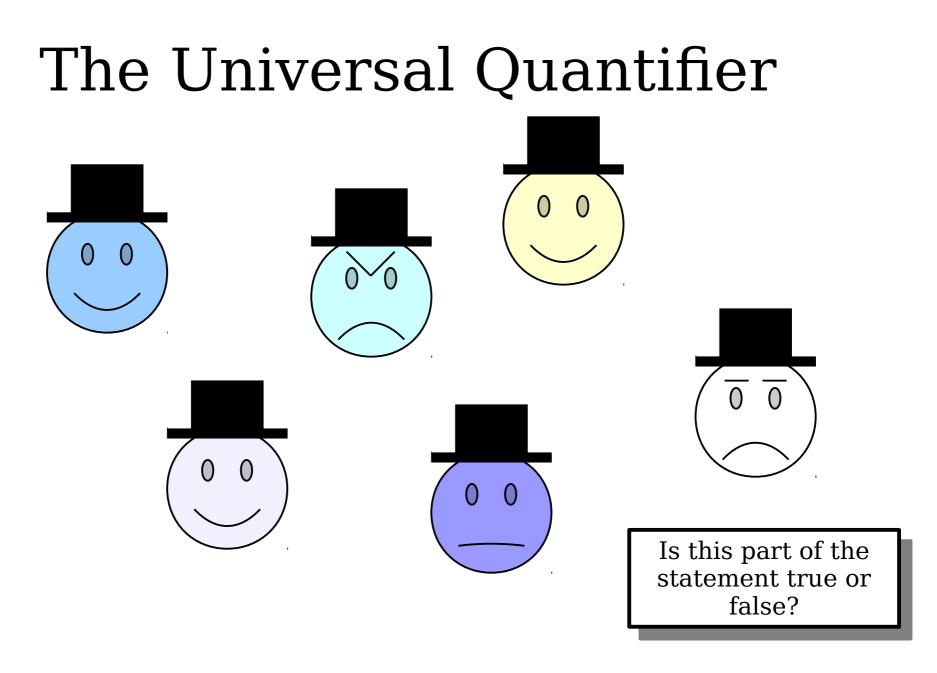


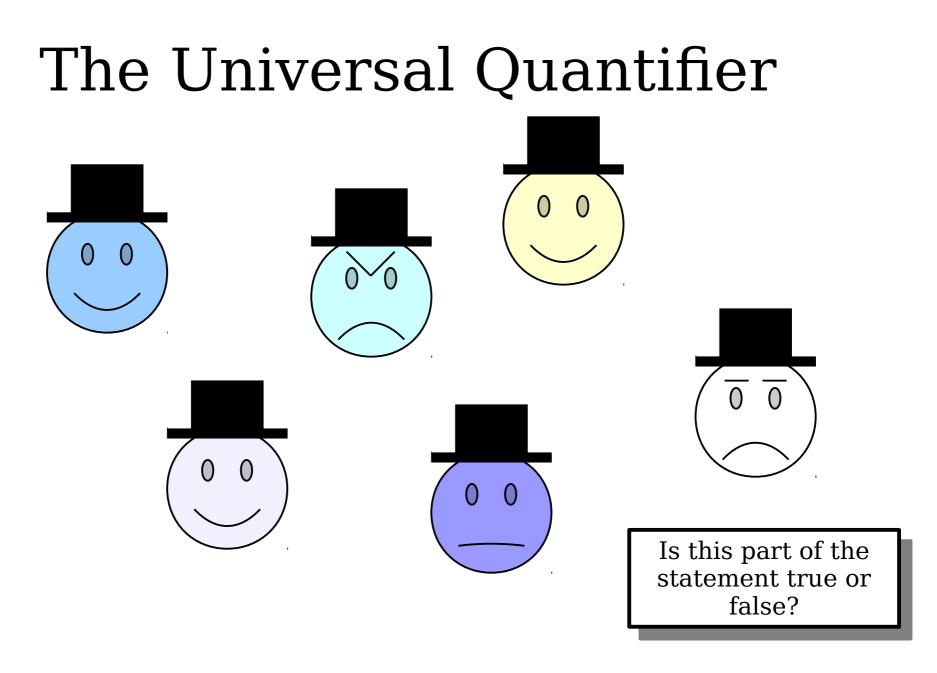
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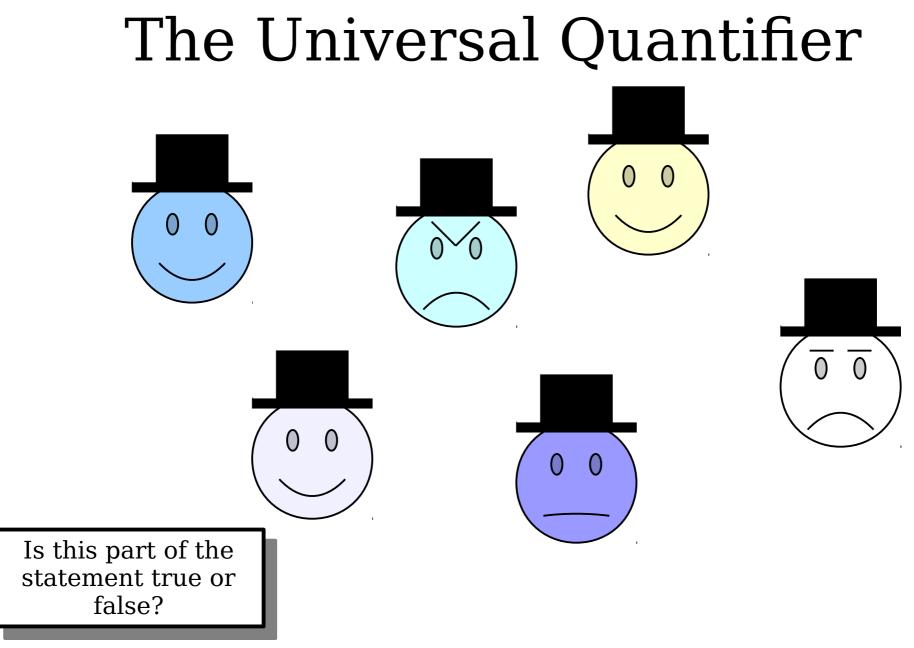
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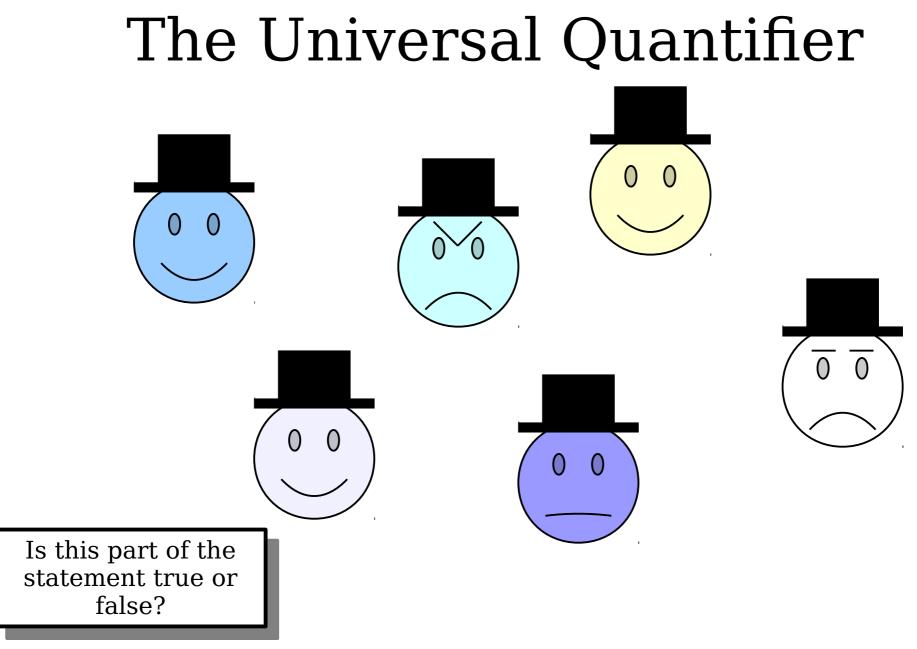


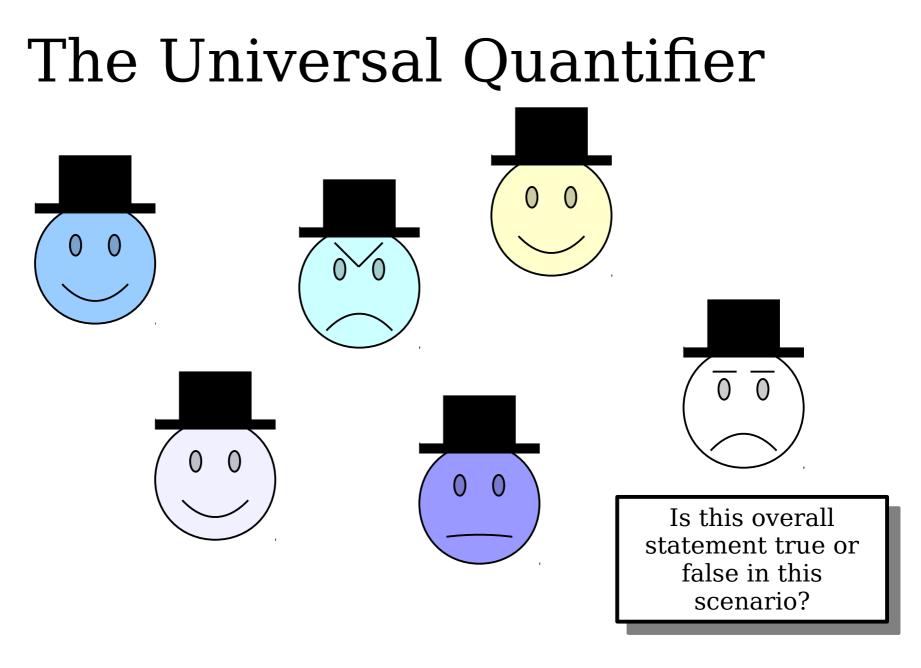


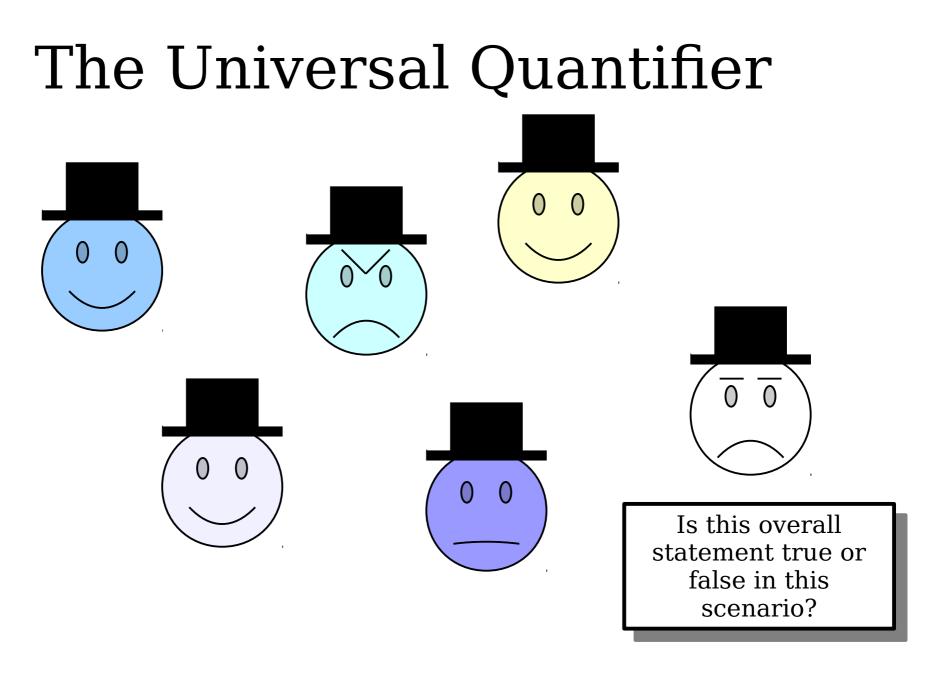












Fun with Edge Cases

 $\forall x. Smiling(x)$

Fun with Edge Cases

Universally-quantified statements are said to be *vacuously true* in empty worlds.

 $\forall x. Smiling(x)$

Let's take a quick break!

Translating into First-Order Logic

Translating Into Logic

- First-order logic is an excellent tool for manipulating definitions and theorems to learn more about them.
- Need to take a negation? Translate your statement into FOL, negate it, then translate it back.
- Want to prove something by contrapositive? Translate your implication into FOL, take the contrapositive, then translate it back.

Translating Into Logic

• When translating from English into firstorder logic, we recommend that you

think of first-order logic as a mathematical programming language.

• Your goal is to learn how to combine basic concepts (quantifiers, connectives, etc.) together in ways that say what you mean.

Using the predicates

- Smiling(x), which states that x is smiling, and
- *WearingHat*(*x*), which states that *x* is wearing a hat,

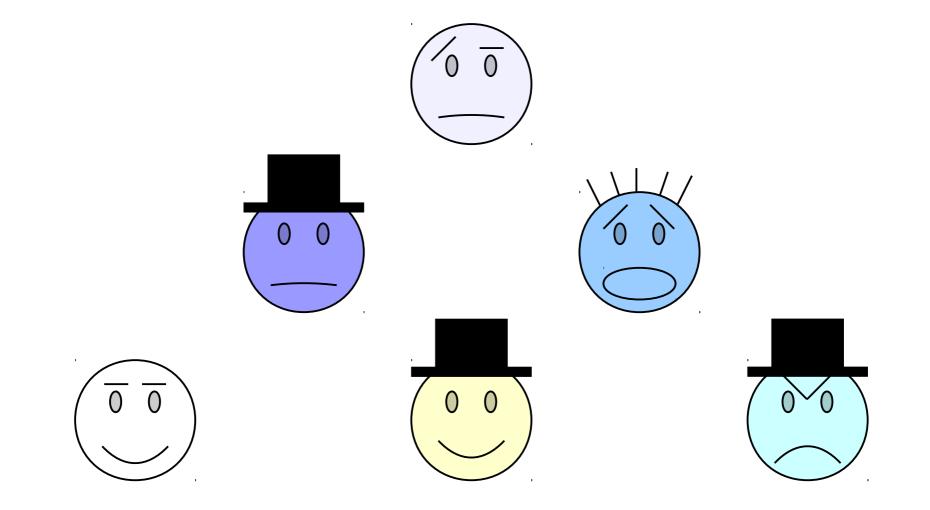
write a sentence in first-order logic that says

some smiling person wears a hat.

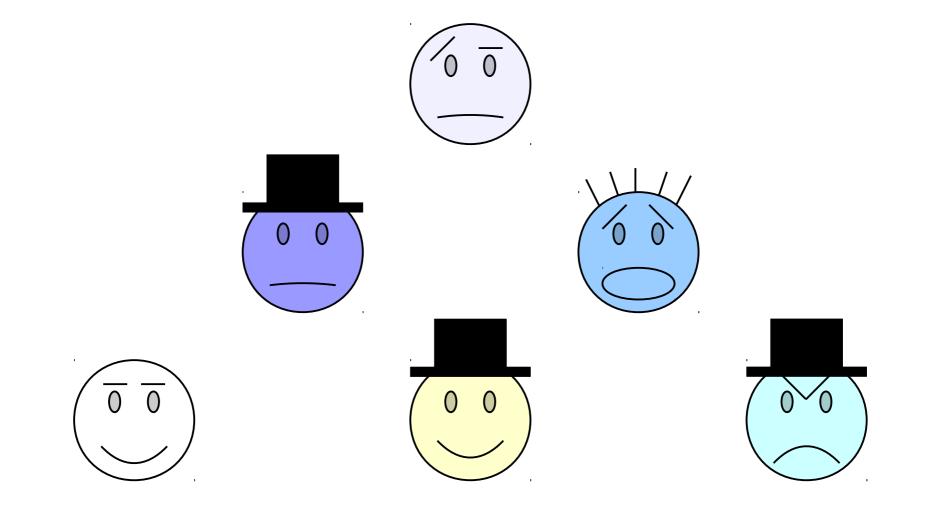
Try it yourself: Give this your best shot – it's okay if you're not sure!

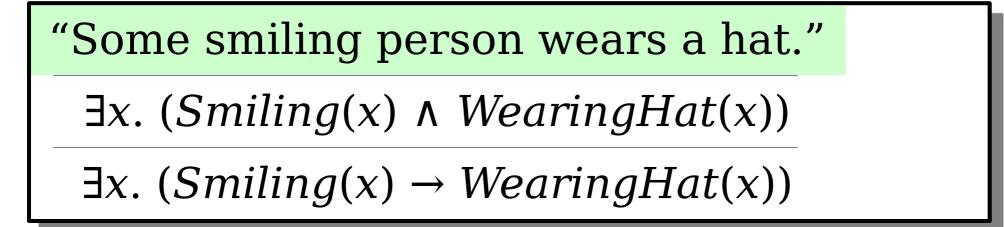
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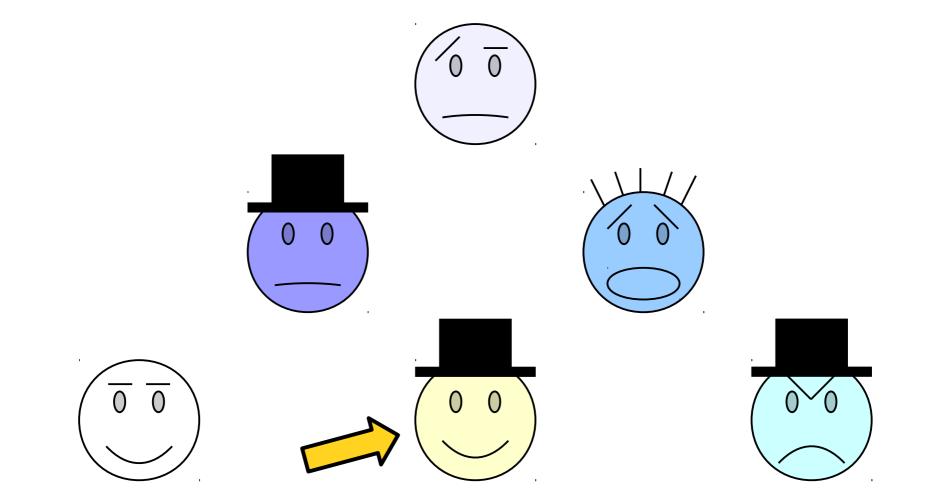
"Some smiling person wears a hat." $\exists x. (Smiling(x) \land WearingHat(x))$ $\exists x. (Smiling(x) \rightarrow WearingHat(x))$

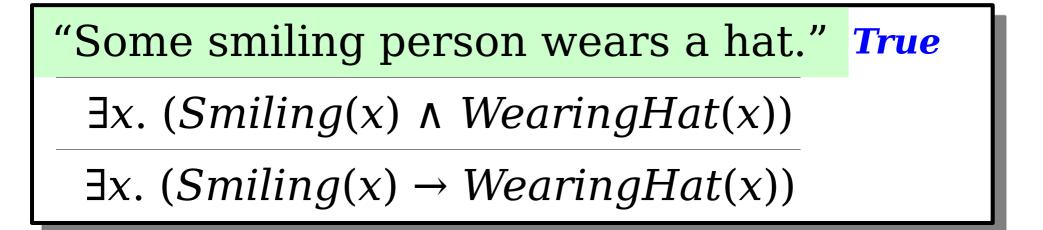


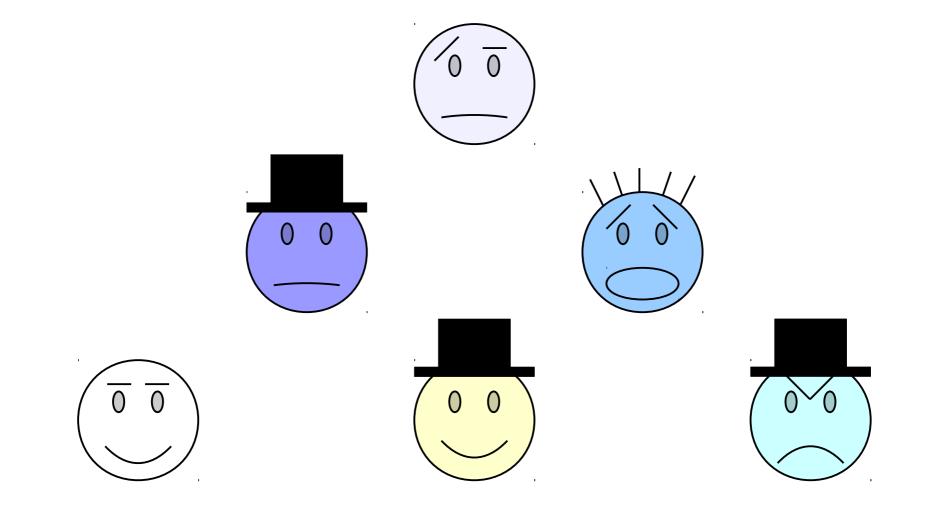
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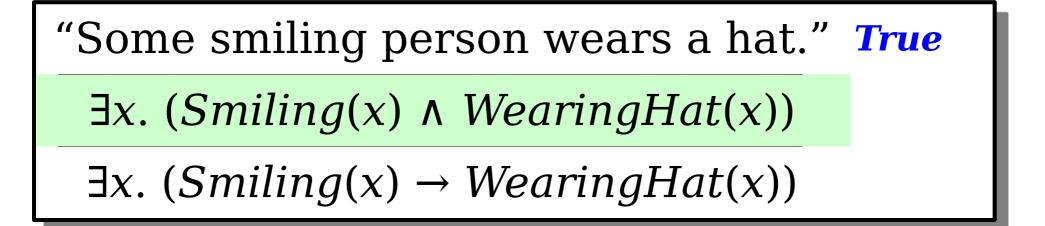


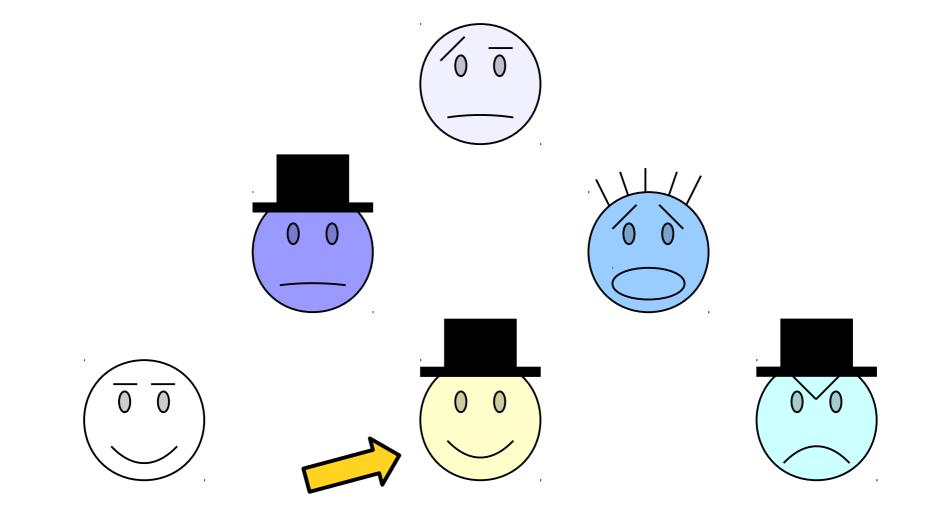


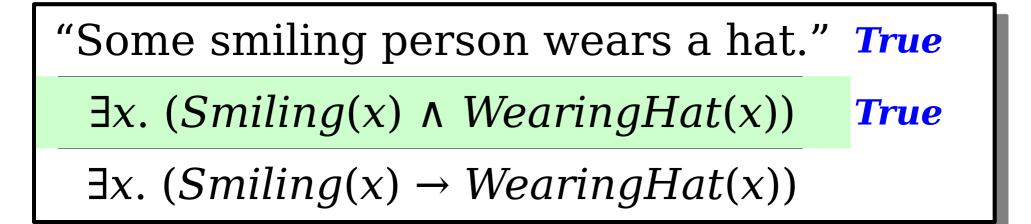


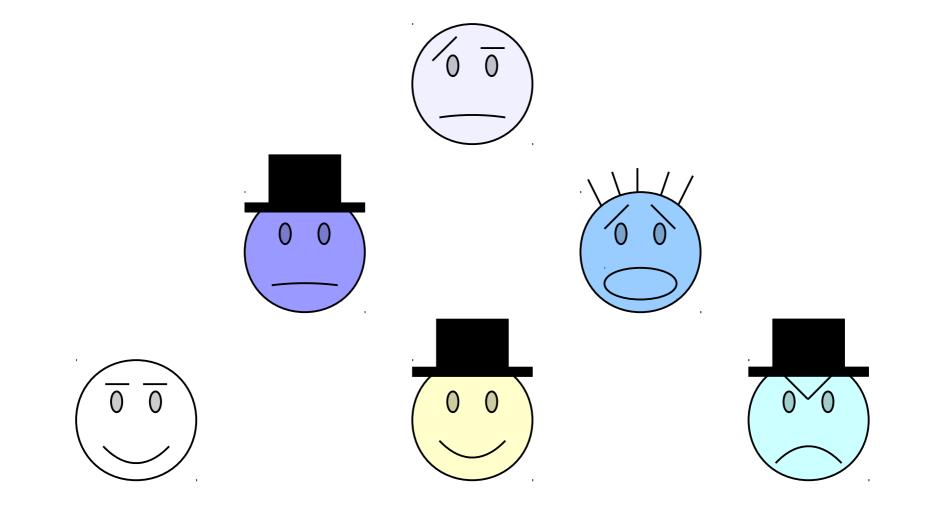


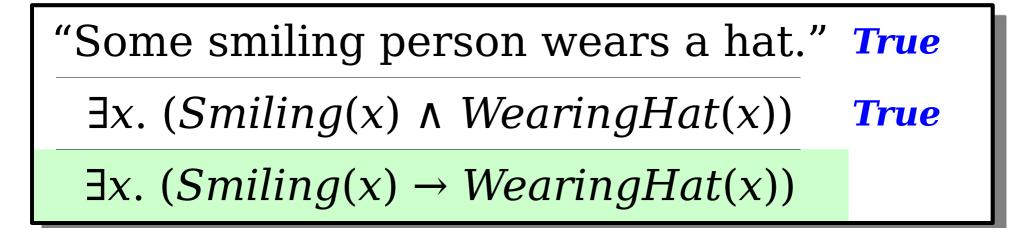


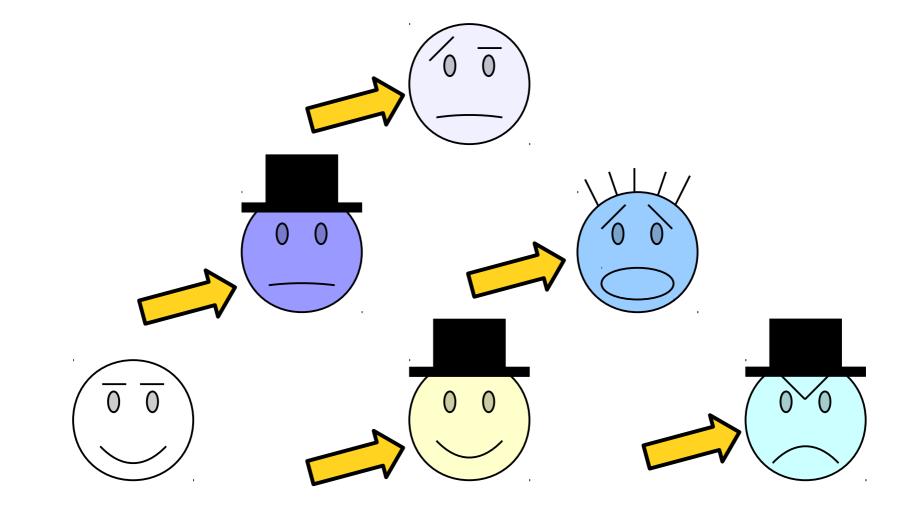


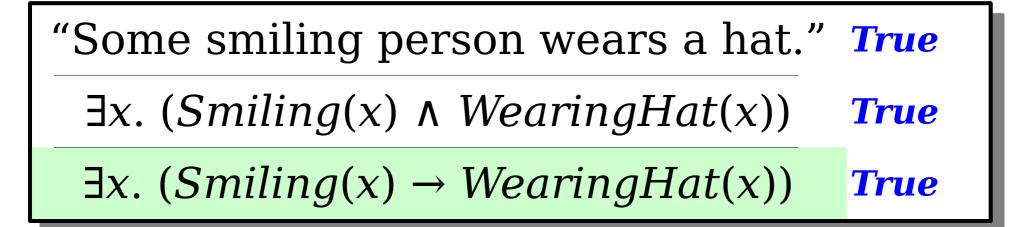


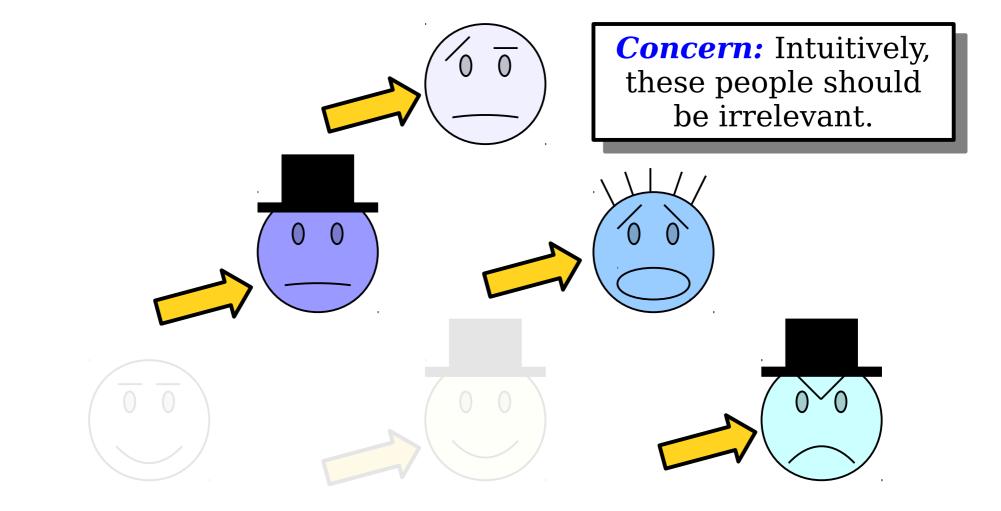


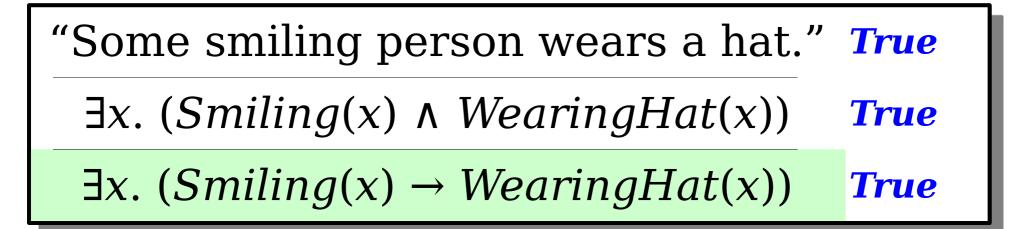




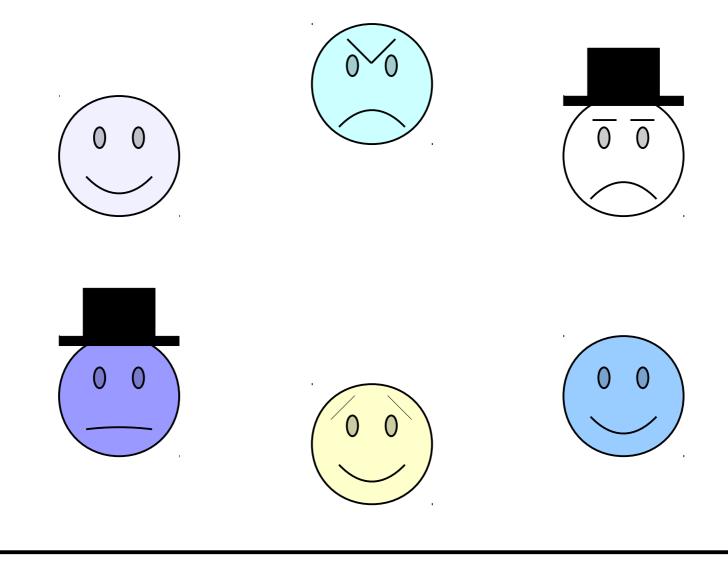


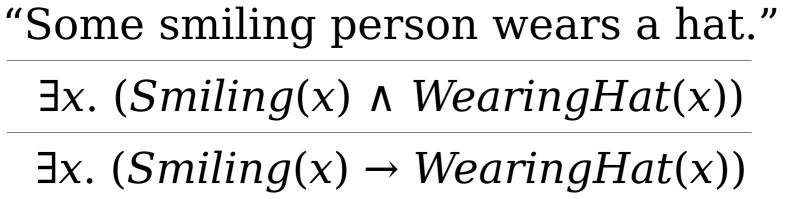


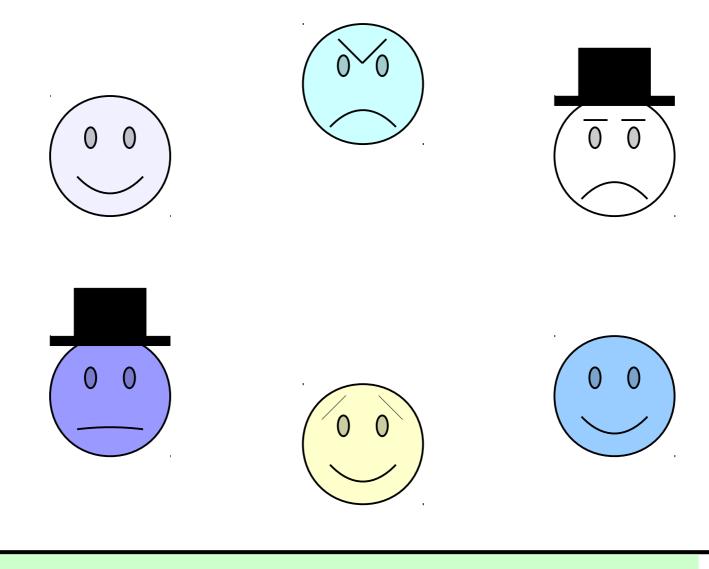




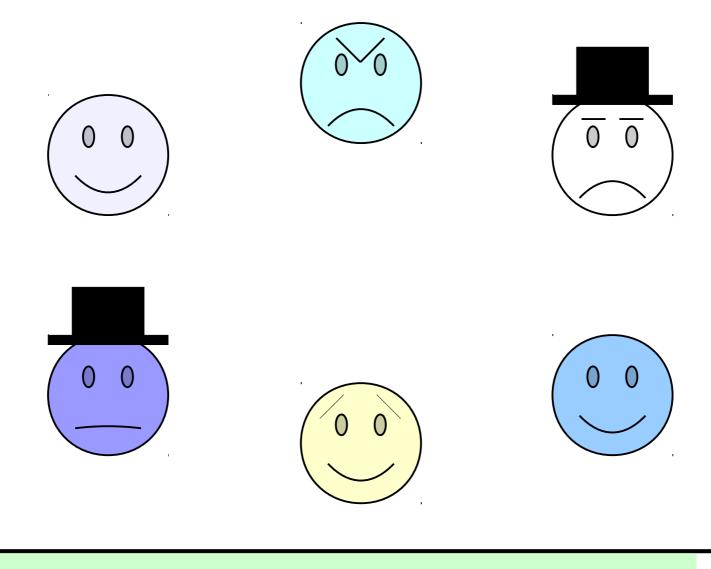
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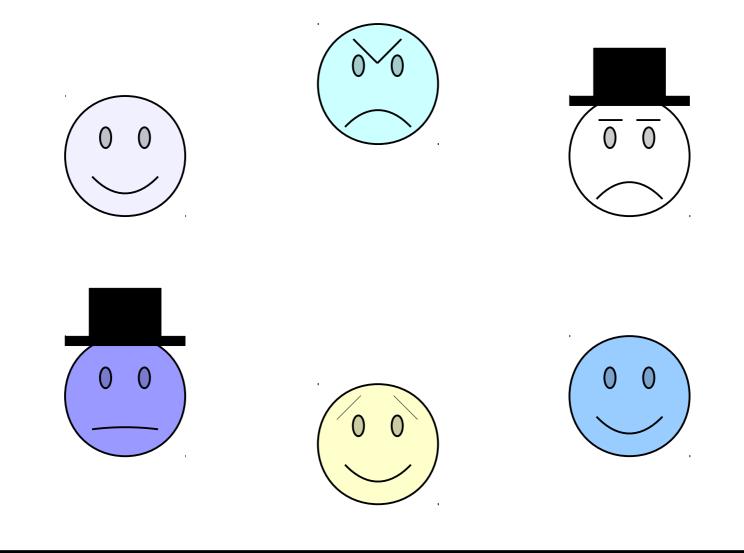


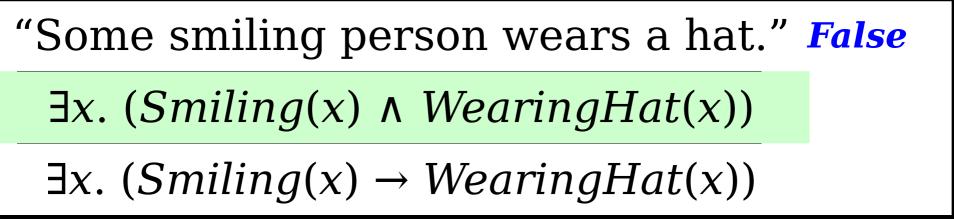


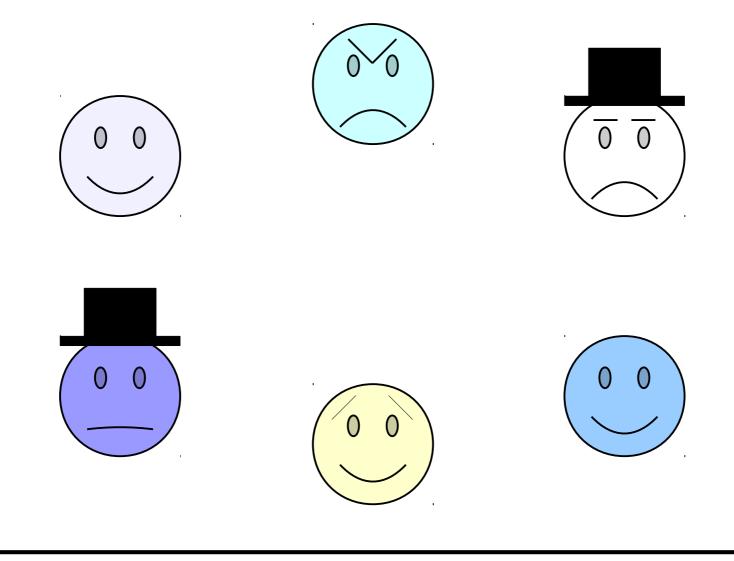
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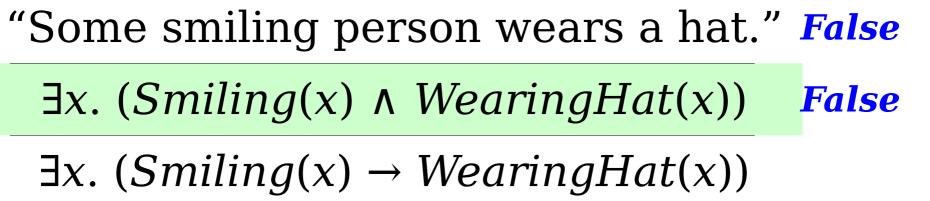


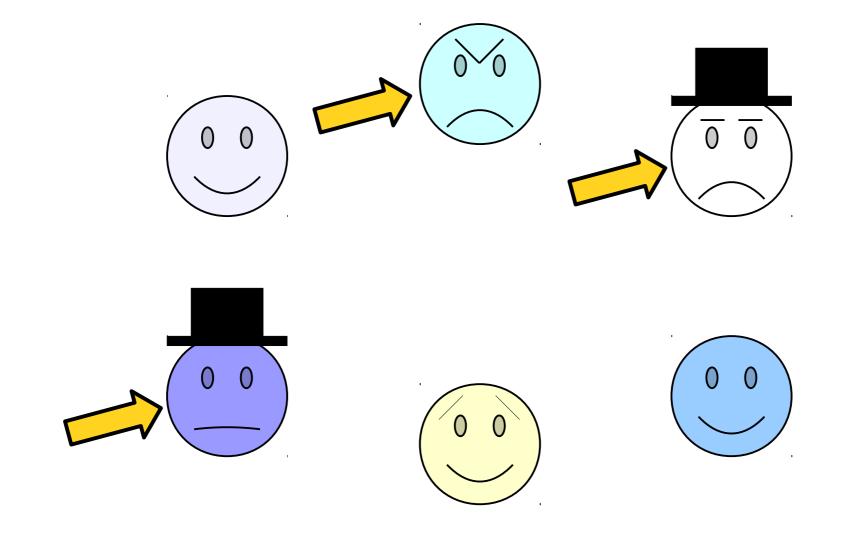
"Some smiling person wears a hat." False $\exists x. (Smiling(x) \land WearingHat(x))$ $\exists x. (Smiling(x) \rightarrow WearingHat(x))$

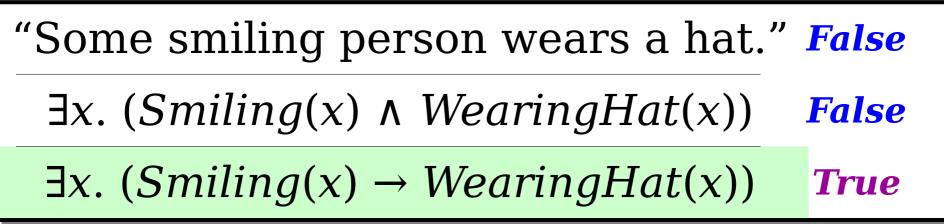


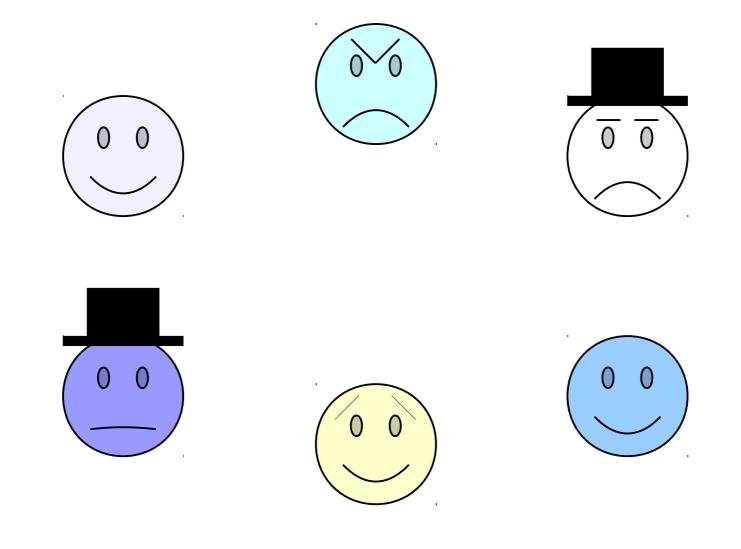


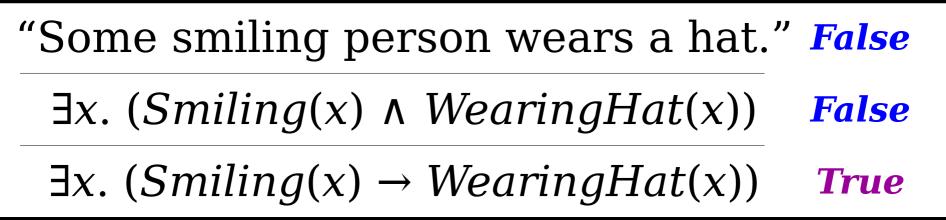


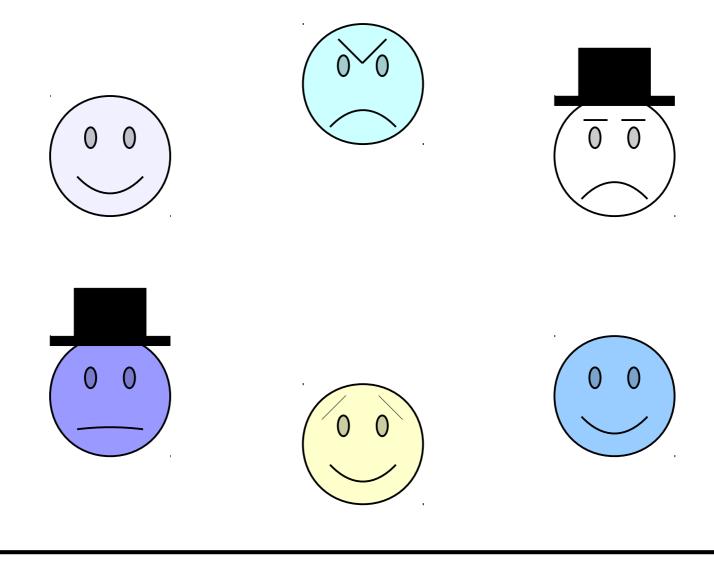


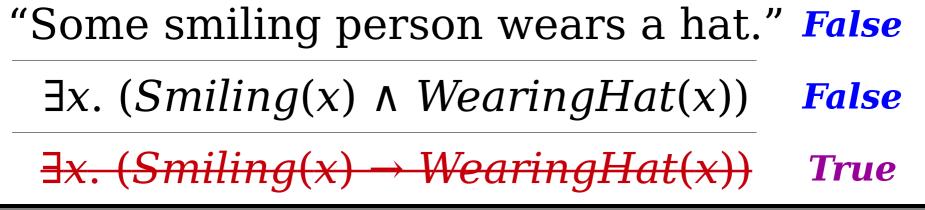












"Some P is a Q"

translates as

 $\exists x. (P(x) \land Q(x))$

Useful Intuition:

Existentially-quantified statements are false unless there's a positive example.

$\exists x. (P(x) \land Q(x))$

If *x* is an example, it *must* have property *P* on top of property *Q*.

Using the predicates

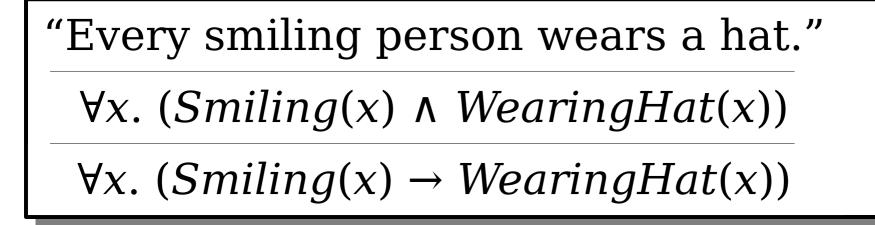
- Smiling(x), which states that x is smiling, and
- *WearingHat*(*x*), which states that *x* is wearing a hat,

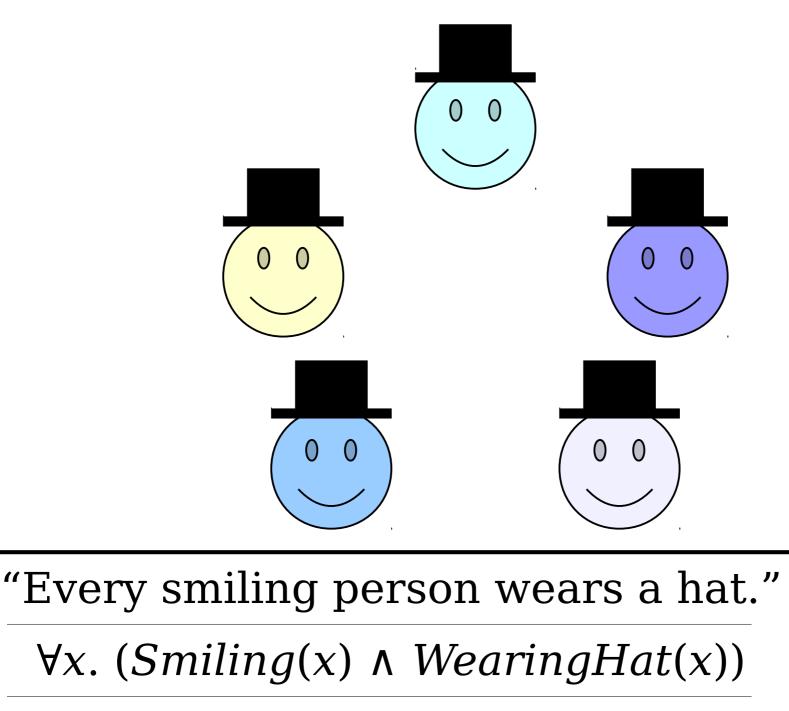
write a sentence in first-order logic that says

every smiling person wears a hat.

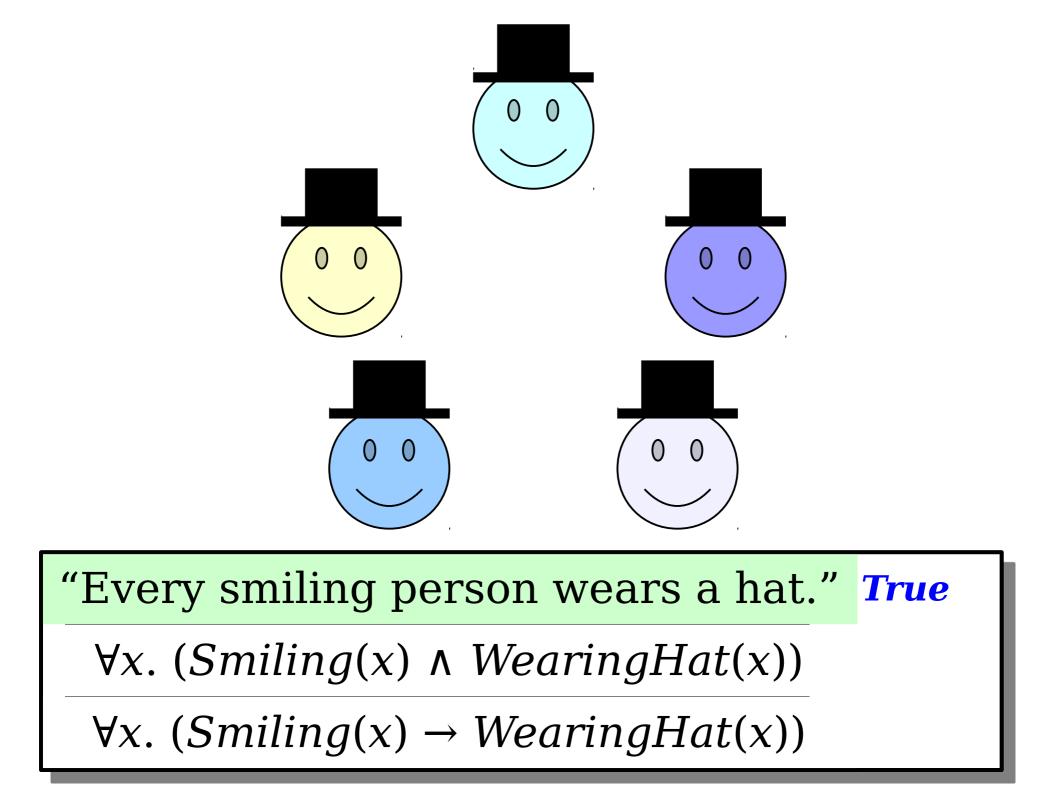
Try it yourself: Give this your best shot – it's okay if you're not sure!

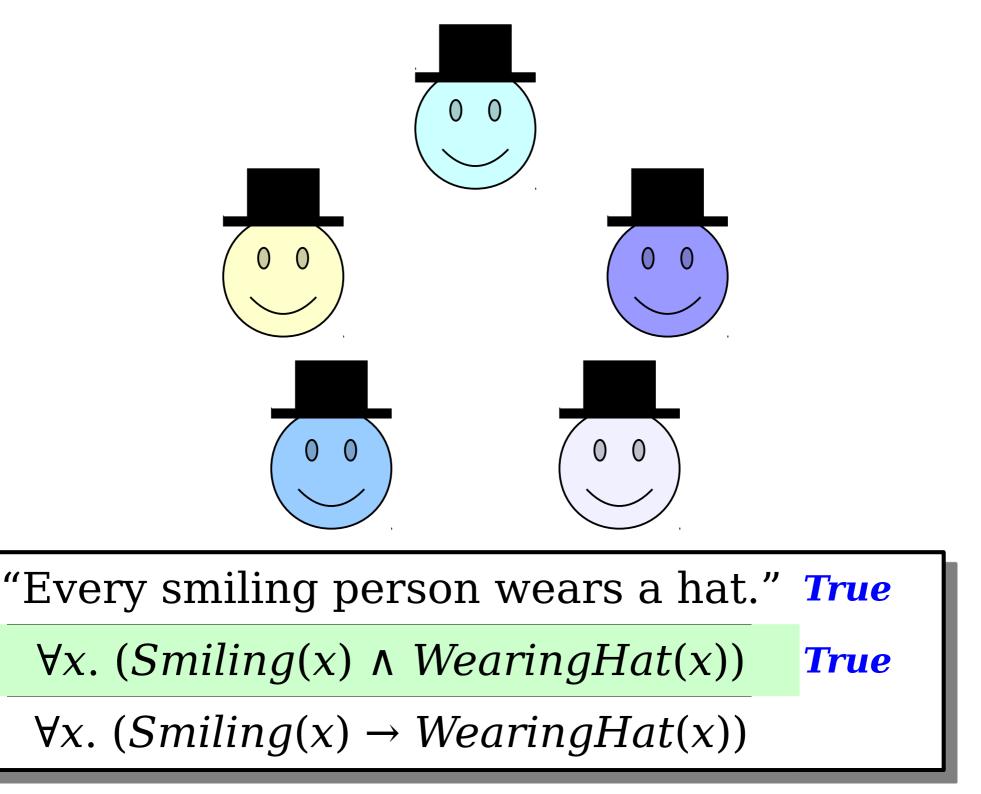
Respond at pollev.com/zhenglian740

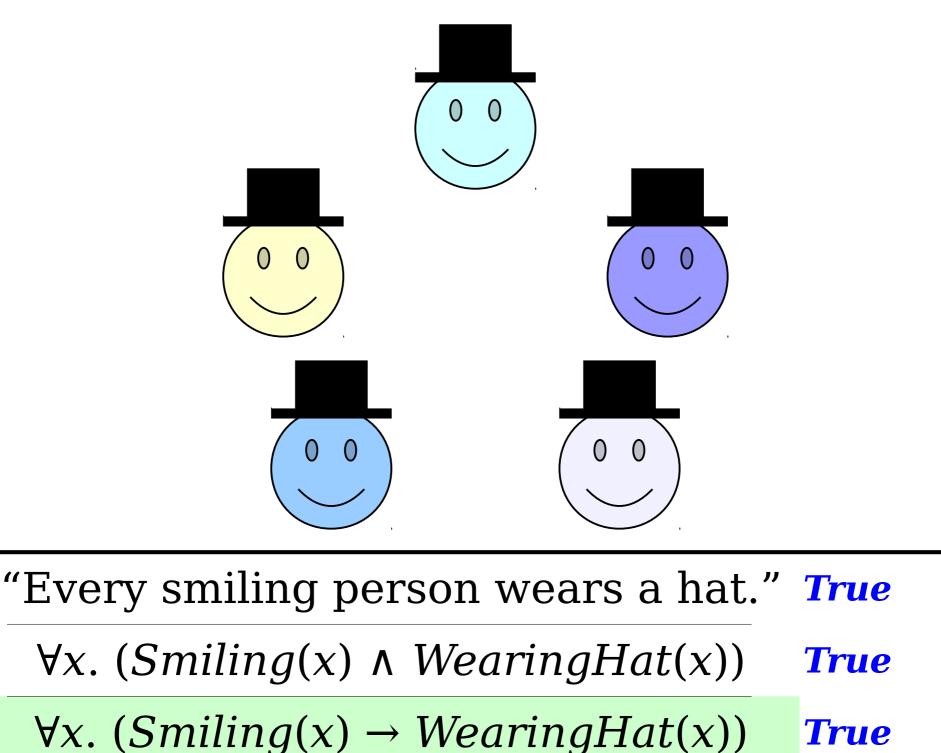


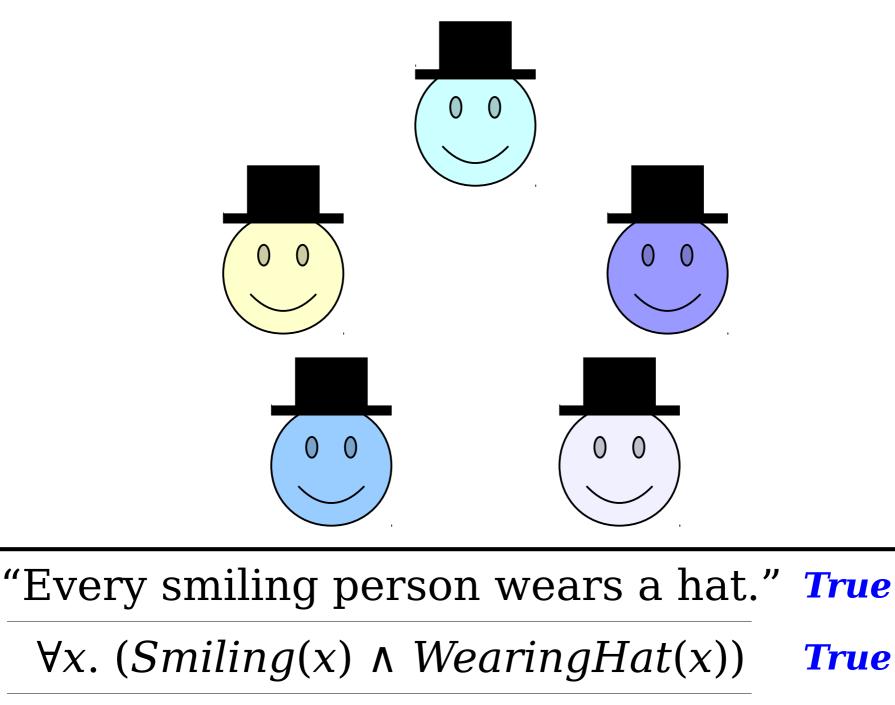


 $\forall x. (Smiling(x) \rightarrow WearingHat(x))$

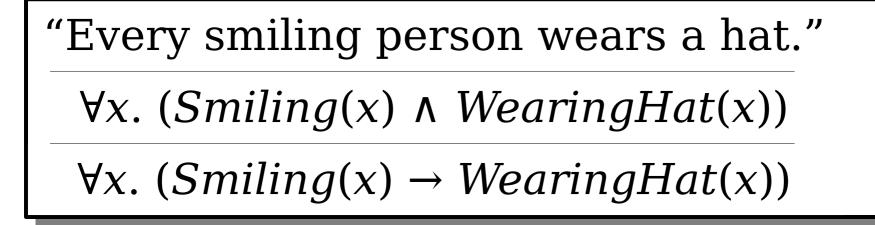


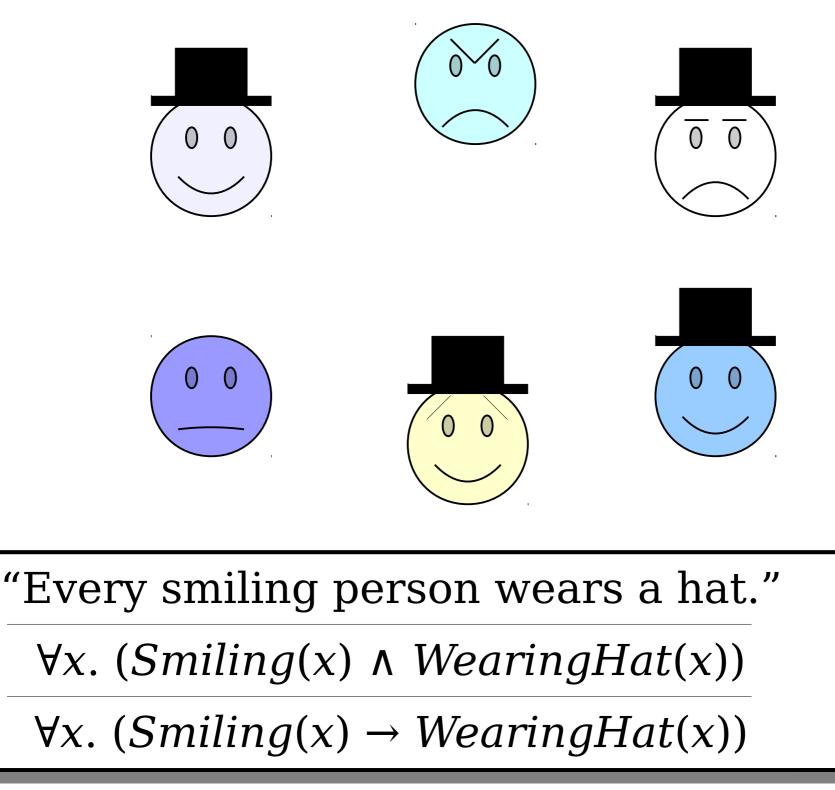


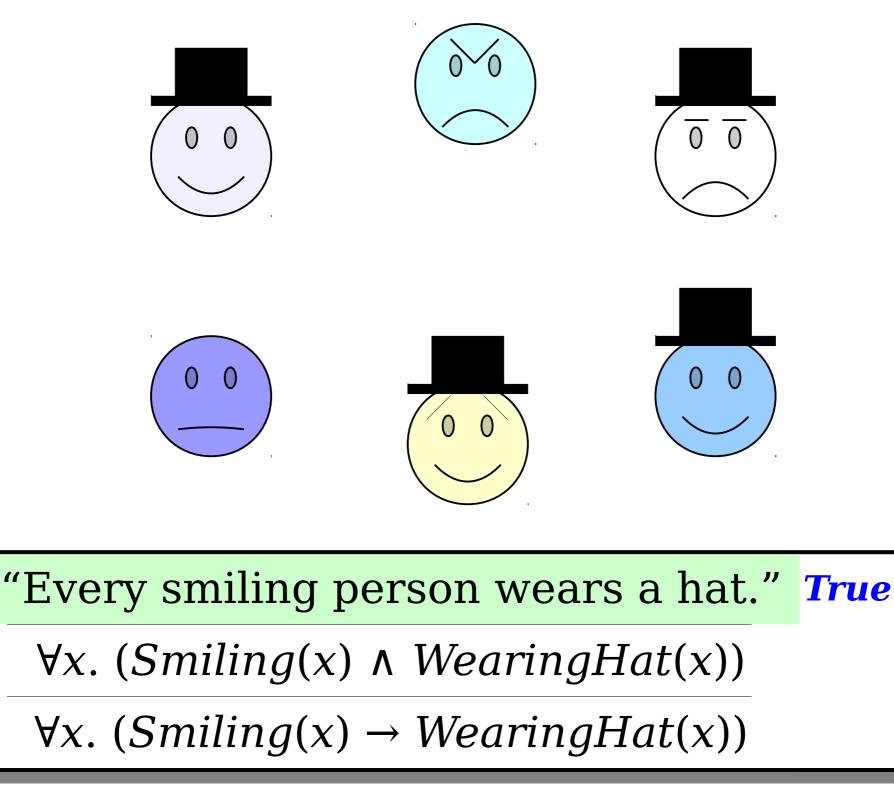


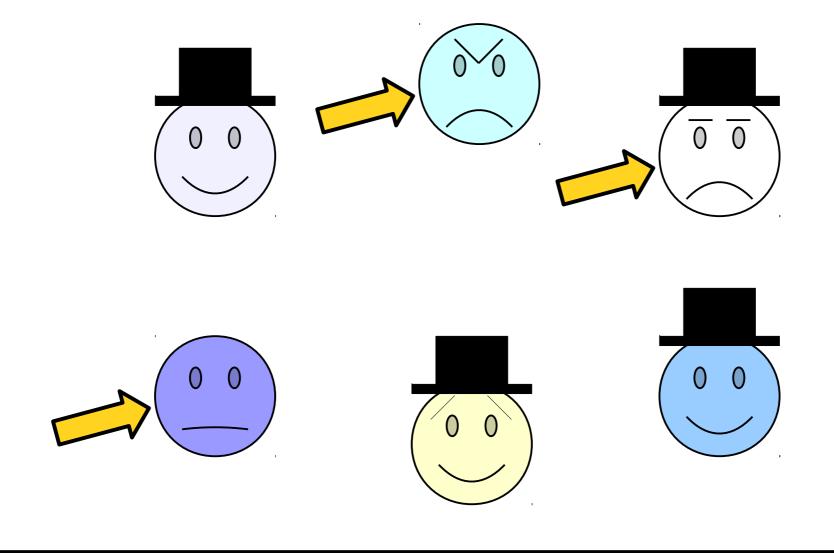


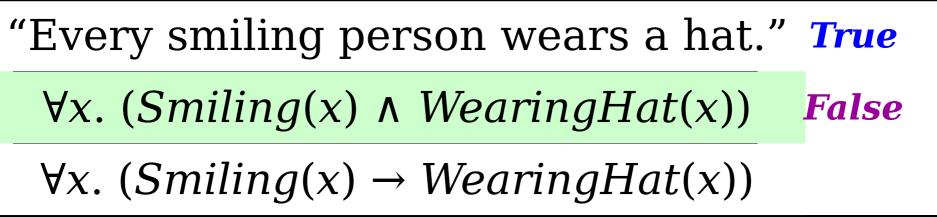
 $\forall x. (Smiling(x) \rightarrow WearingHat(x))$ **True**

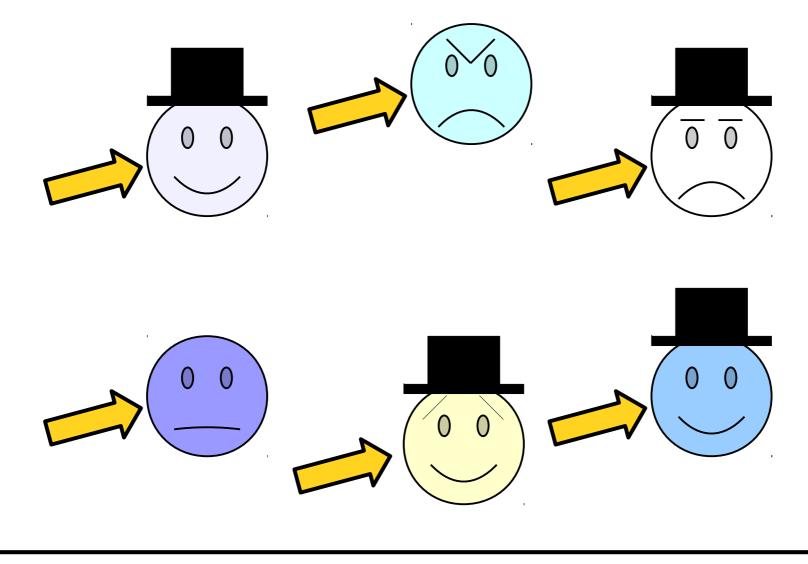


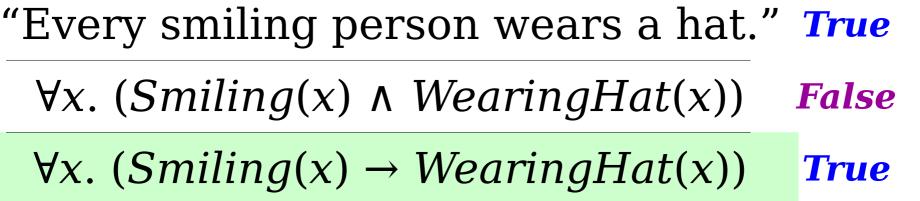


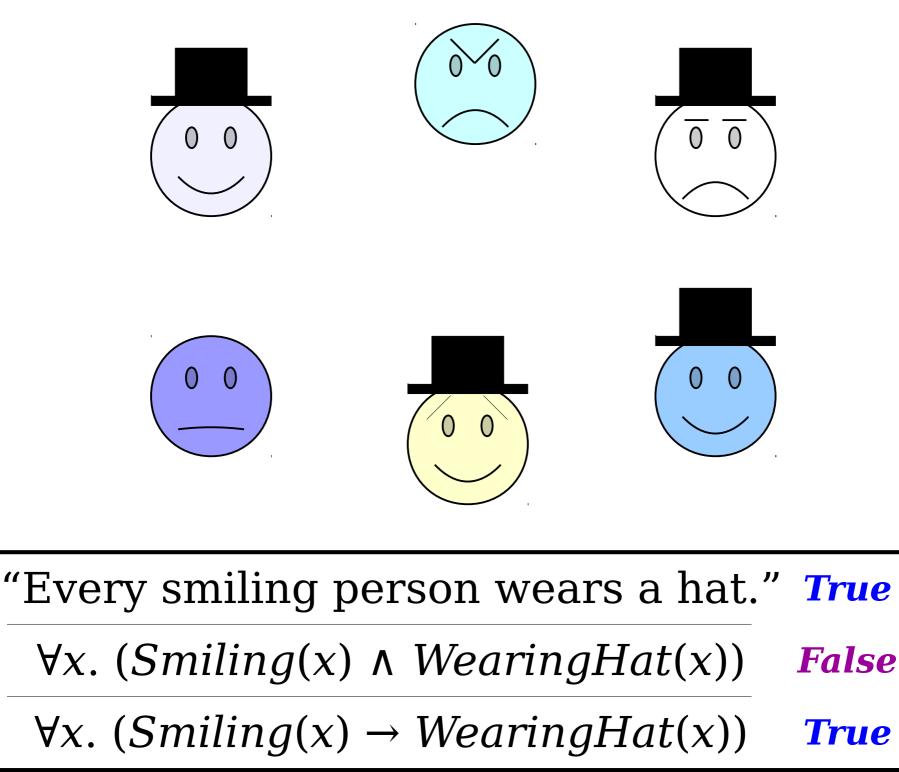


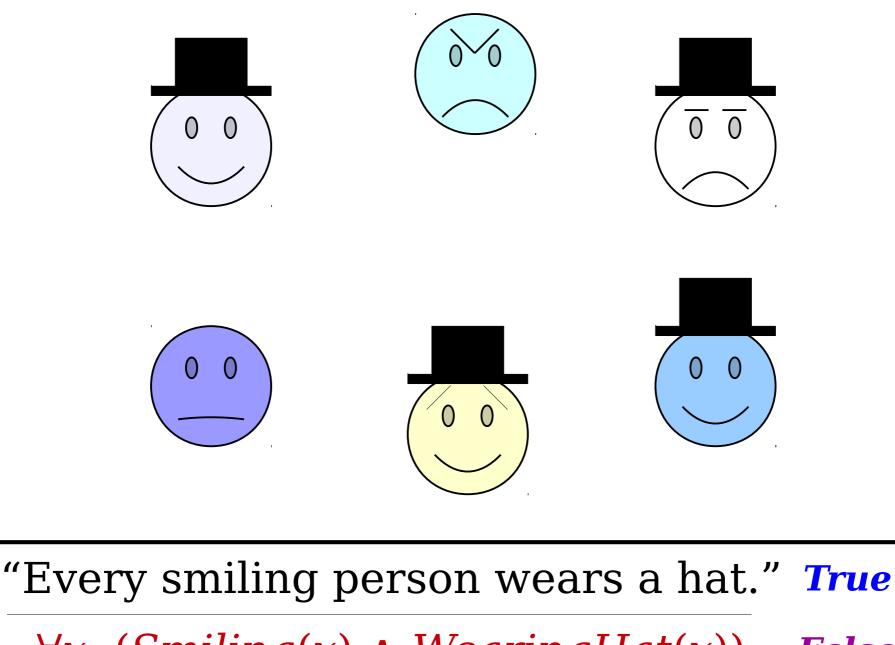














"All P's are Q's"

translates as

 $\forall x. \ (P(x) \rightarrow Q(x))$

Useful Intuition:

Universally-quantified statements are true unless there's a counterexample.

$\forall x. \ (P(x) \rightarrow Q(x))$

If *x* is a counterexample, it *must* have property *P* but not have property *Q*.

Good Pairings

• The \forall quantifier *usually* is paired with \rightarrow .

 $\forall x. \ (P(x) \rightarrow Q(x))$

• The \exists quantifier *usually* is paired with \land .

$\exists x. (P(x) \land Q(x))$

- In the case of ∀, the → connective prevents the statement from being *false* when speaking about some object you don't care about.
- In the case of \exists , the \land connective prevents the statement from being *true* when speaking about some object you don't care about.

Proofwriting Workshop

An Incorrect Set Theory Proof

▲ **Incorrect!** ▲ **Proof:** Consider arbitrary sets A, B, and C where $C \subseteq A \cup B$.

This means that every element of *C* is in either *A* or *B*. If all elements of *C* are in *A*, then $C \subseteq A$. Alternately, if everything in *C* is in *B*, then $C \subseteq B$. In either case, everything inside of *C* has to be contained in at least one of these sets, so the theorem is true.

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This is just repeating definitions and not making specific claims about specific variables.

▲ **Incorrect!** ▲ **Proof:** Consider arbitrary sets A, B, and C where $C \subseteq A \cup B$.

This means that every element of *C* is in either *A* or *B*. If all elements of *C* are in *A*, then $C \subseteq A$. Alternately, if everything in *C* is in *B*, then $C \subseteq B$. In either case, everything inside of *C* has to be contained in at least one of these sets, so the theorem is true.

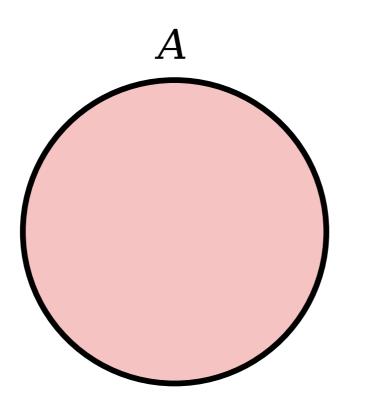
Why is this bad?

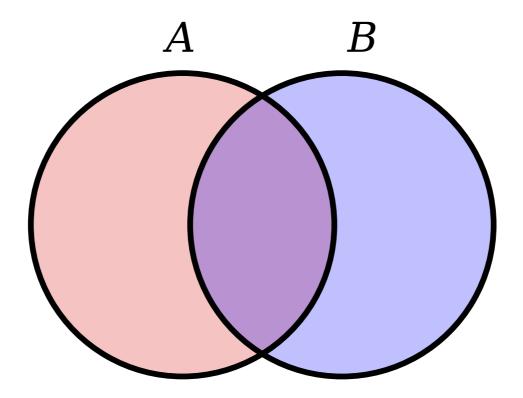
▲ **Incorrect!** ▲ **Proof:** Consider arbitrary sets A, B, and C where $C \subseteq A \cup B$.

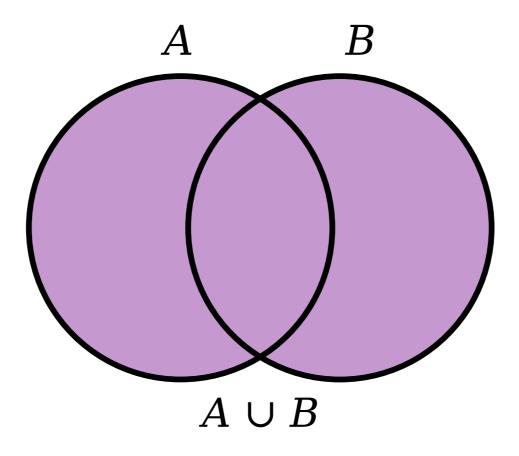
This means that every element of C is in either A or B. If all elements of C are in A, then $C \subseteq A$. Alternately, if everything in C is in B, then $C \subseteq B$. In either case, everything inside of C has to be contained in at least one of these sets, so the

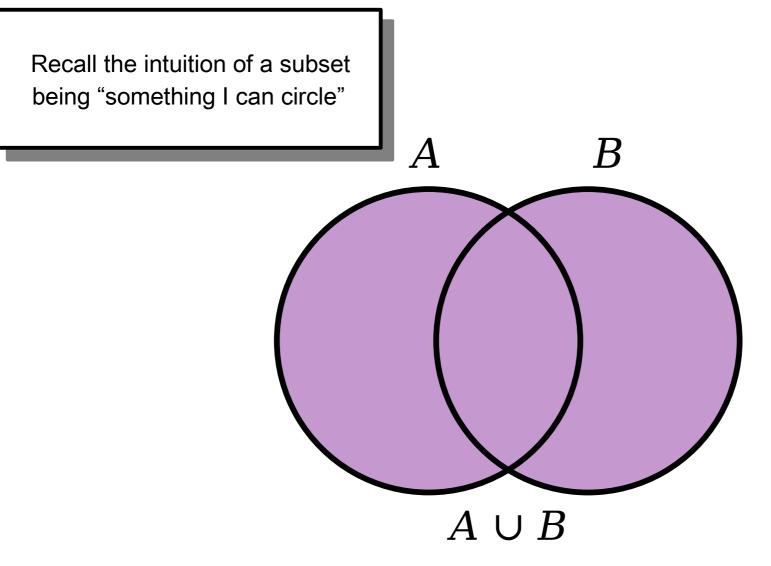
theorem is true.

While this claim is true, it does not imply the theorem is true. In fact, this theorem is actually false.



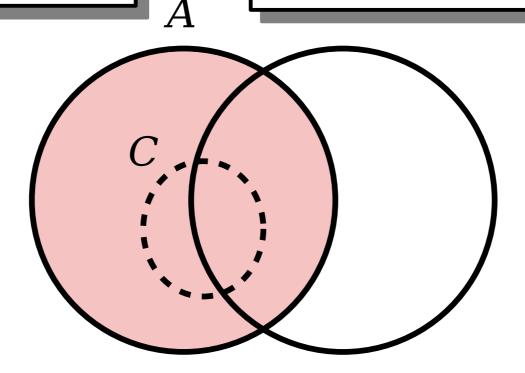


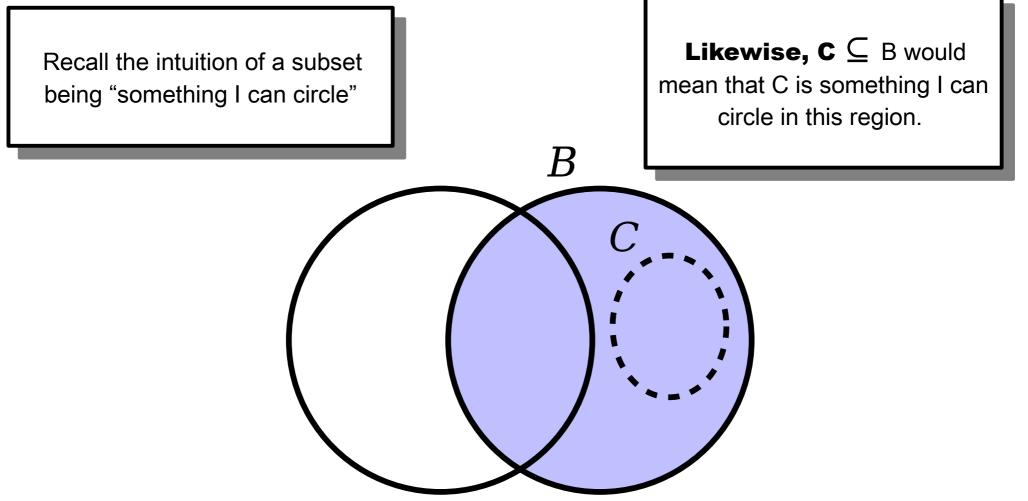


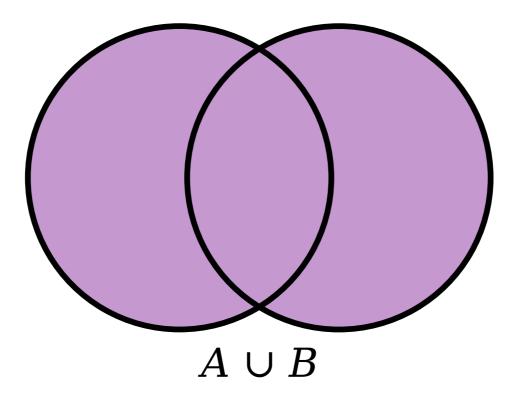


Recall the intuition of a subset being "something I can circle"

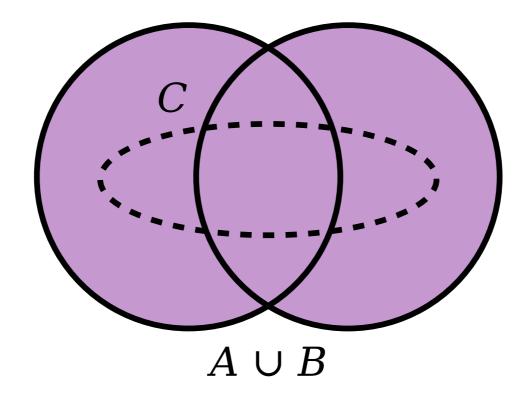
So C ⊆ A would mean that C is something I can circle in this region.







But when I look at $A \cup B$, I can draw C as a circle containing elements from both A and B!



But when I look at $A \cup B$, I can draw C as a circle containing elements from both A and B!

Do you see why this circle is in neither A nor B?

 $A \cup B$

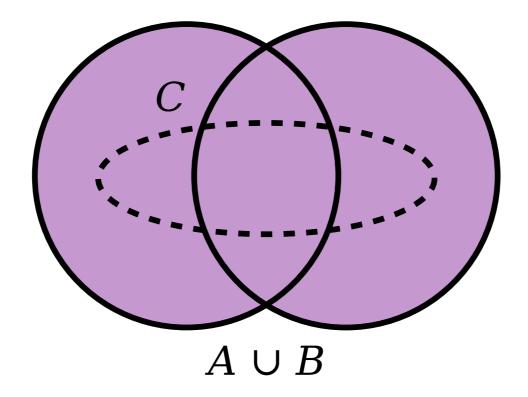
Let's Draw Some Pictures!

Claim: If A, B, and C are sets and $C \subseteq A \cup B$, then $C \subseteq A$

or $C \subseteq B$.

Using this visual intuition, come up with a choice of sets *A*, *B*, and *C* that show this claim is false.

Respond at pollev.com/zhenglian740

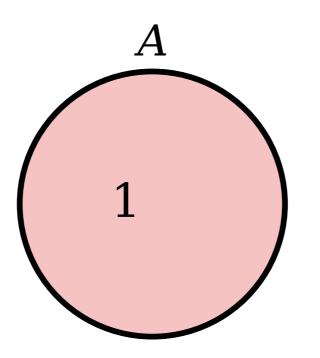


Proofs vs. Disproofs

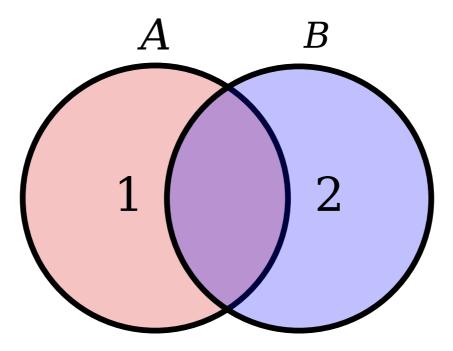
- A **proof** is an argument that explains why some **theorem** is true.
- A *disproof* is an argument that explains why some *claim* is false.
- You've seen lots of examples of proofs. What does a disproof look like?

Disproof: We will show that there are sets A, B, and C where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$.

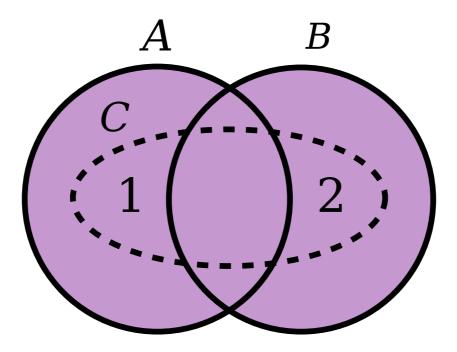
Disproof: We will show that there are sets *A*, *B*, and *C* where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}$



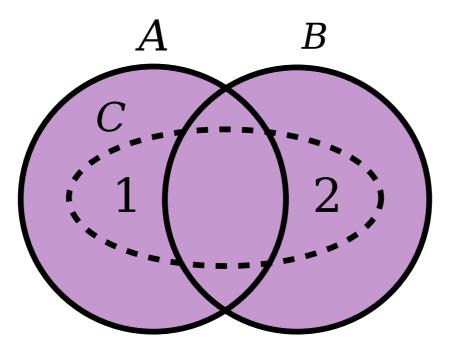
Disproof: We will show that there are sets A, B, and C where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}, B = \{2\}$



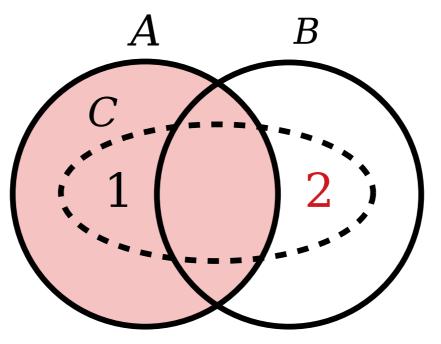
Disproof: We will show that there are sets A, B, and C where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}, B = \{2\}$, and $C = \{1, 2\}$.



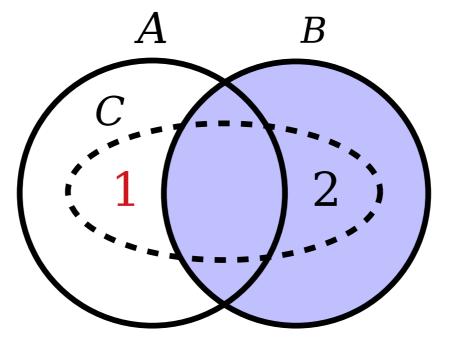
Disproof: We will show that there are sets A, B, and C where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}, B = \{2\}$, and $C = \{1, 2\}$. Now notice that $\{1, 2\} \subseteq A \cup B$ so $C \subseteq A \cup B$



Disproof: We will show that there are sets *A*, *B*, and *C* where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}, B = \{2\}$, and $C = \{1, 2\}$. Now notice that $\{1, 2\} \subseteq A \cup B$ so $C \subseteq A \cup B$, but $C \not\subseteq A$ because $2 \in C$ but $2 \notin A$



Disproof: We will show that there are sets *A*, *B*, and *C* where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}, B = \{2\}$, and $C = \{1, 2\}$. Now notice that $\{1, 2\} \subseteq A \cup B$ so $C \subseteq A \cup B$, but $C \not\subseteq A$ because $2 \in C$ but $2 \notin A$, and $C \not\subseteq B$ because $1 \in C$ but $1 \notin B$.



Disproof: We will show that there are sets *A*, *B*, and *C* where $C \subseteq A \cup B$, but $C \not\subseteq A$ and $C \not\subseteq B$. Consider the sets $A = \{1\}, B = \{2\}$, and $C = \{1, 2\}$. Now notice that $\{1, 2\} \subseteq A \cup B$ so $C \subseteq A \cup B$, but $C \not\subseteq A$ because $2 \in C$ but $2 \notin A$, and $C \not\subseteq B$ because $1 \in C$ but $1 \notin B$.

Thus we've found a set *C* which is a subset of $A \cup B$ but is not a subset of either *A* or *B*, which is what we needed to show.

Proofwriting Advice

- Be *very wary* of proofs that speak generally about "all objects" of a particular type.
 - As you've just seen, it's easy to accidentally prove a false statement at this level of detail.
 - Making broad, high-level claims often indicates deeper logic errors or conceptual misunderstanding (like *code smell* but for proofs!)

Proofwriting Advice

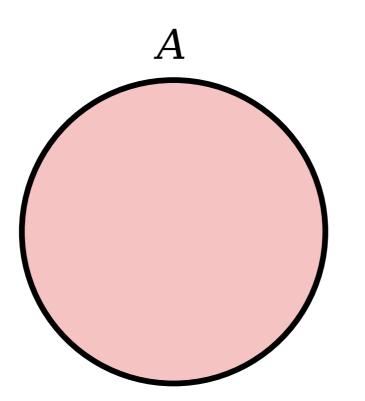
A Very Good Idea: After you've written a draft of a proof, run through all of the points on the Proofwriting Checklist.

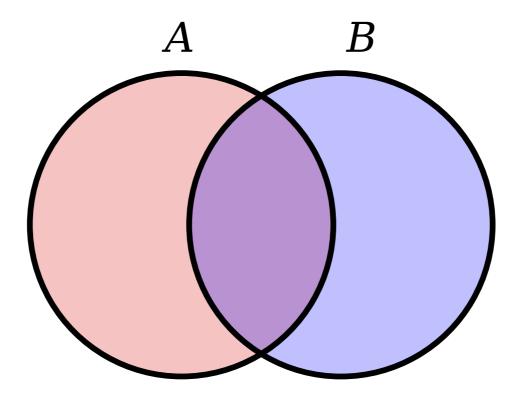
• This is a *great* exercise that you can do with a partner!

Proofs on Subsets

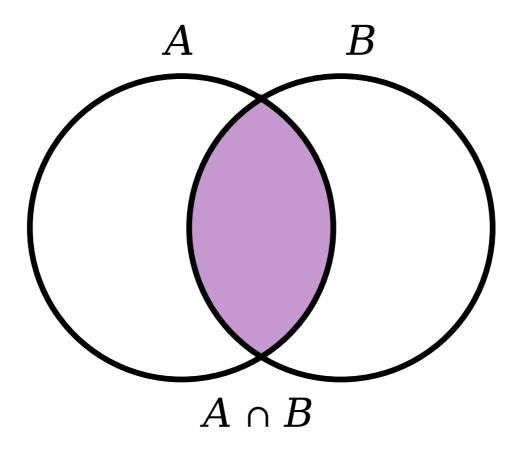
Hold on, isn't this the claim we just disproved?

Notice that that's an intersection, not a union! It turns out that this claim is actually true.

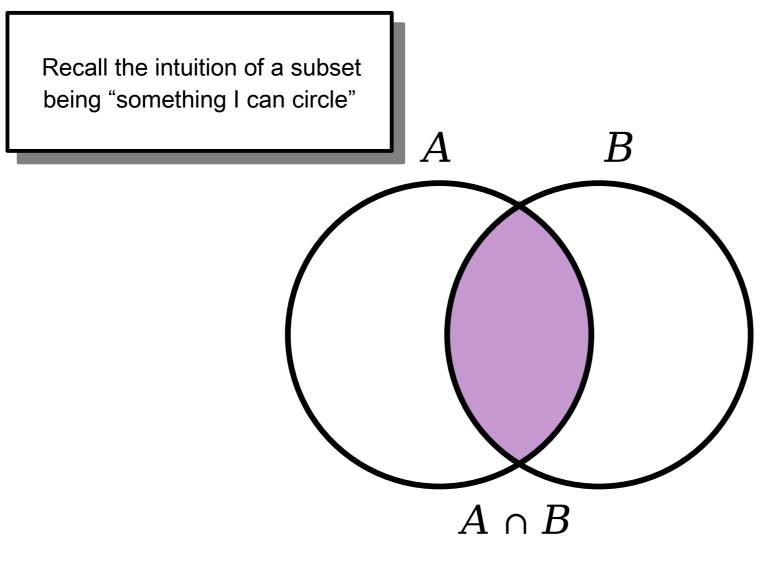


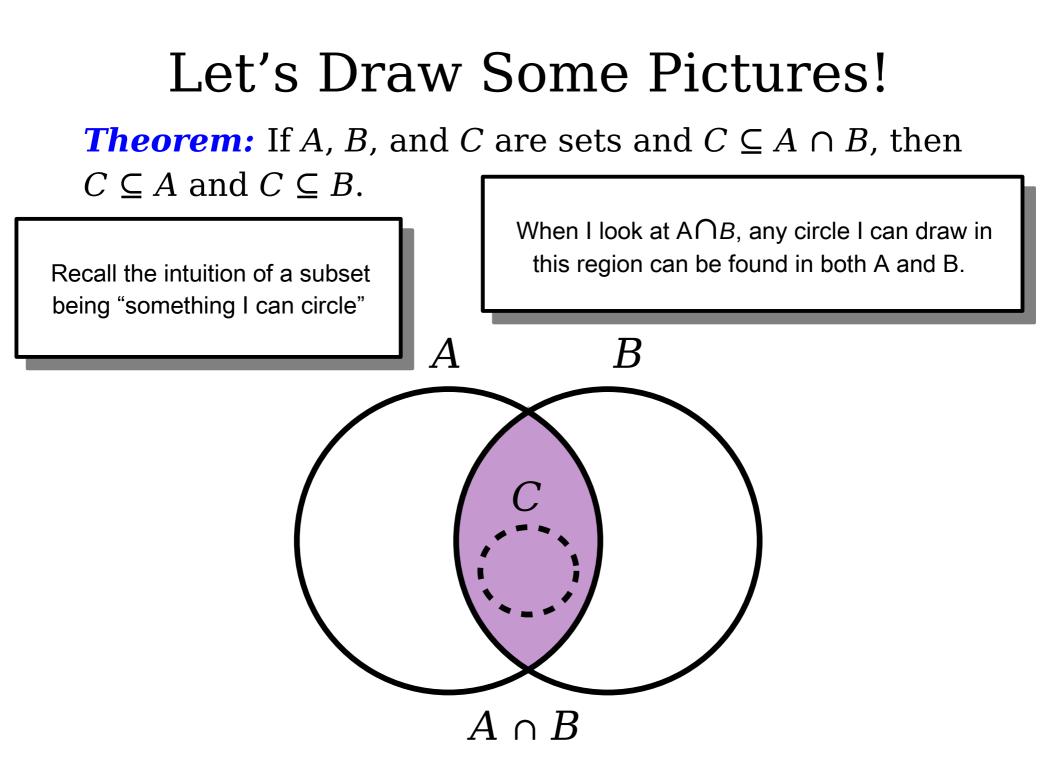


Let's Draw Some Pictures! **Theorem:** If A, B, and C are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.



Let's Draw Some Pictures! **Theorem:** If A, B, and C are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.





Let's Draw Some Pictures! Theorem: If A, B, and C are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

 $A \cap B$

A

This is a great visual intuition to see why the theorem is true. Now we have to drill down to the level of individual elements to write the proof.

What We're Assuming

What We Need To Show

When confronted with a theorem to prove, the first step is to make sure you understand where you're starting and where you're going.

What We're Assuming

- A, B, and C are sets
- C⊆A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

A great proofwriting strategy is to **write down relevant definitions**. This gives you a better sense of what you

need to prove and what tools you have at hand.

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

Before we start:

- What is the definition of subset?
- How do you prove that one set is a subset of another?
- If you know that one set is a subset of can you conclude?

another, what

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

Definition: If S and T are sets, then $S \subseteq T$ when for every $x \in S$, we have $x \in T$.

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

Definition: If *S* and *T* are sets, then $S \subseteq T$ when for every $x \in S$, we have $x \in T$.

To prove that $S \subseteq T$: Pick an arbitrary $x \in S$, then prove $x \in T$.

If you know that $S \subseteq T$: If you have an $x \in S$, you can conclude $x \in T$.

What We're Assuming

- A, B, and C are sets
- $C \subseteq A \cap B$

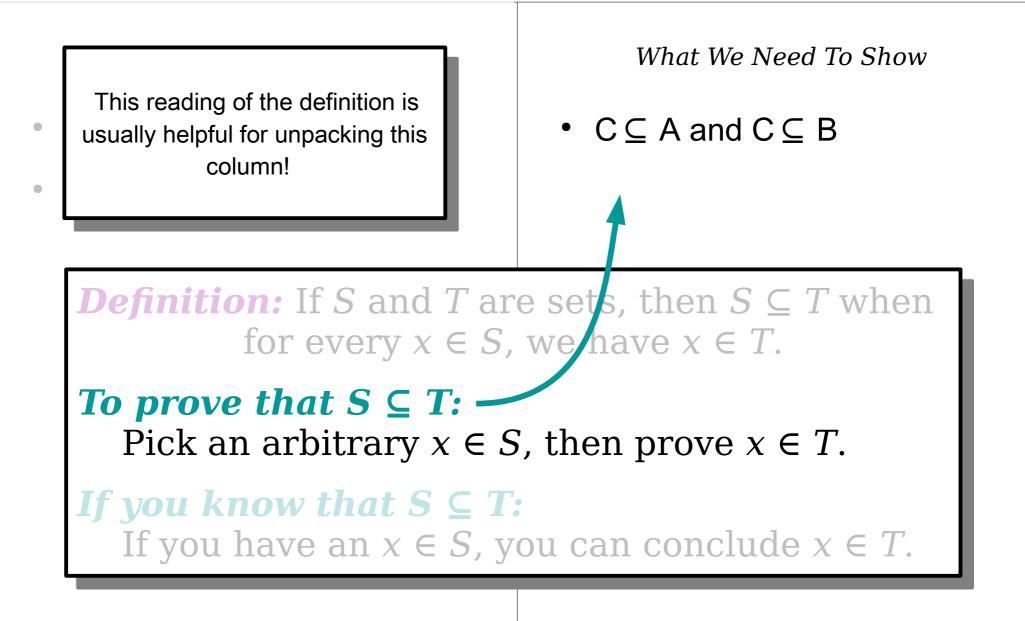
What We Need To Show

•
$$C \subseteq A$$
 and $C \subseteq B$

Definition: If S and T are sets, then $S \subseteq T$ when for every $x \in S$, we have $x \in T$.

To prove that $S \subseteq T$: Pick an arbitrary $x \in S$, then prove $x \in T$.

If you know that $S \subseteq T$: If you have an $x \in S$, you can conclude $x \in T$.



What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

Our Tools

• In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$ What We Need To Show

•
$$C \subseteq A$$
 and $C \subseteq B$

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

Our Tools

• In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$ How can we apply this general template to our specific problem?

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

Our Tools

• In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$ What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
 - Pick an $x \in C$, show that $x \in A$
 - Pick an $x \in C$, show that $x \in B$

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

Our Tools

• In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$ What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
 - Pick an $x \in C$, show that $x \in A$
 - Pick an $x \in C$, show that $x \in B$

Now we know that ultimately, we're going to have to do these two things. Let's see what tools we have that can get us here!

What We're Assuming

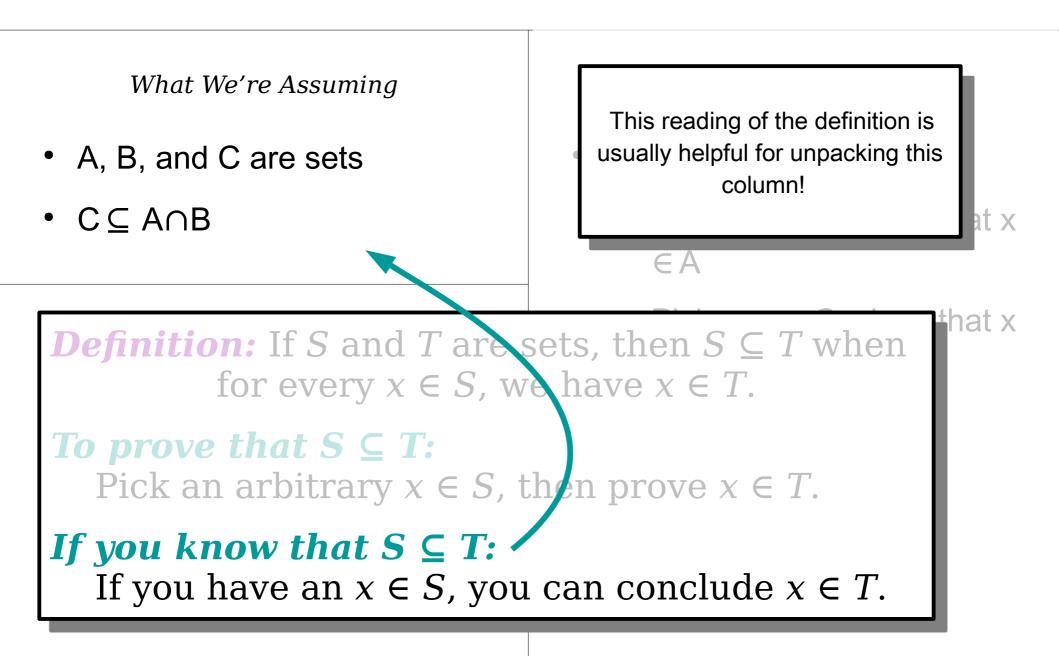
- A, B, and C are sets
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Our Tools

• In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
 - Pick an $x \in C$, show that $x \in A$
 - Pick an $x \in C$, show that $x \in B$



What We're Assuming

• A, B, and C are sets

• C ⊆ A∩B

Our Tools

• In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$ What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

- Pick an $x \in C$, show that $x \in A$
- Pick an $x \in C$, show that $x \in B$

Before we continue:

- What is the definition of set intersection?

What We're Assuming

• A, B, and C are sets

• C ⊆ A∩B

What We Need To Show

• $C \subseteq A$ and $C \subseteq B$

• Pick an $x \in C$, show that $x \in A$

Definition: The set $S \cap T$ is the set where, for any x: $x \in S \cap T$ when $x \in S$ and $x \in T$

What We're Assuming

• A, B, and C are sets

• C ⊆ A∩B

What We Need To Show

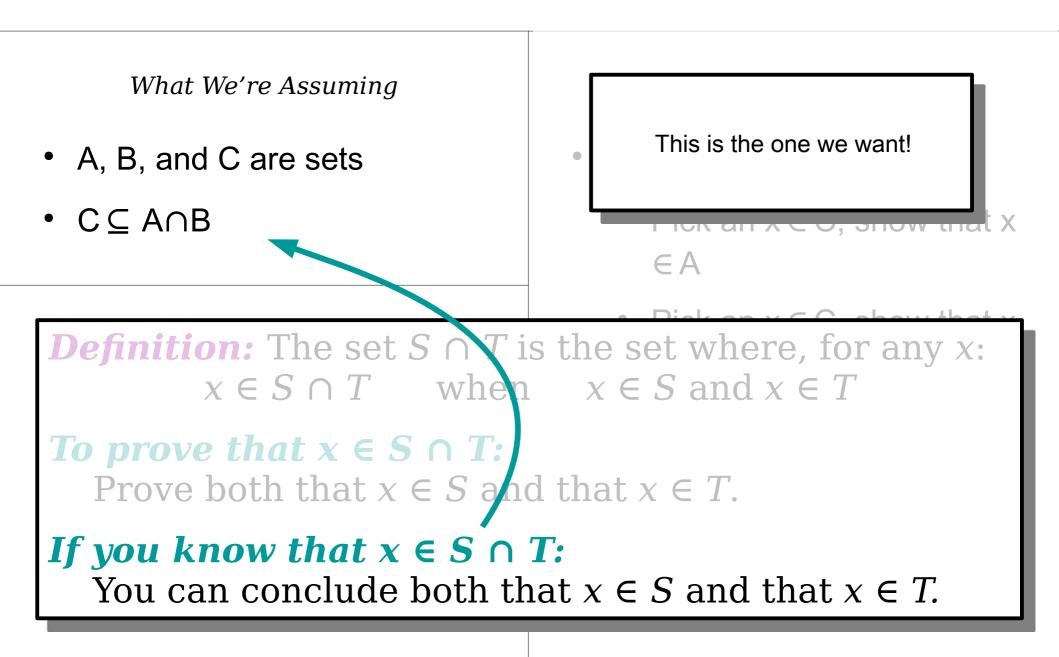
• $C \subseteq A$ and $C \subseteq B$

• Pick an $x \in C$, show that $x \in A$

Definition: The set $S \cap T$ is the set where, for any x: $x \in S \cap T$ when $x \in S$ and $x \in T$

To prove that $x \in S \cap T$: Prove both that $x \in S$ and that $x \in T$.

If you know that $x \in S \cap T$: You can conclude both that $x \in S$ and that $x \in T$.



What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$.
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
 - Pick an $x \in C$, show that $x \in A$
 - Pick an $x \in C$, show that $x \in B$

What We're Assuming

- A, B, and C are sets
- C ⊆ A∩B

Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$.
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
 - Pick an $x \in C$, show that $x \in A$
 - Pick an $x \in C$, show that $x \in B$

Let's go and try and do the proof with what we've got here!

Rough Outline

• Assume $C \subseteq A \cap B$

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
 - Pick an $x \in C$

• Conclude $x \in A$

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
 - Pick an $x \in C$

What goes here?

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
 - Pick an $x \in C$
 - x ∈ A ∩ B
 - Conclude $x \in A$

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
 - Pick an $x \in C$
 - $x \in A \cap B$
 - Conclude $x \in A$

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 - Pick an $x \in C$
 - $x \in A \cap B$
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 - Conclude $x \in A$

We also need to prove that $C \subseteq B$.

Notice that if you take the outline here and literally swap the variable A for the variable B, you get a working proof.

Rough Outline

- Assume $C \subseteq B \cap A$
- Proving $C \subseteq B$
 - Pick an $x \in C$
 - $x \in B \cap A$
 - $x \in B$ and $x \in A$
 - Conclude $x \in B$

In a case like this where your proof would have two completely symmetric branches, it's fine to write up just one and say "by symmetry, [the other branch] is also true."

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
 - Pick an $x \in C$
 - $x \in A \cap B$
 - $x \in A$ and $x \in B$
 - Conclude $x \in A$

Try it yourself: Take a few minutes and write up a proof of the theorem using this outline.

Then share your proof with your neighbors and critique each other!

Respond at pollev.com/zhenglian740

Proof: Let A, B, and C be arbitrary sets where $C \subseteq A \cap B$. We need to show that $C \subseteq A$ and $C \subseteq B$. Because the roles of A and B in this proof are symmetric, we can just prove that $C \subseteq A$.

Choose any element $x \in C$. Since $C \subseteq A \cap B$, we know that $x \in A \cap B$. This tells us that $x \in A$ and $x \in B$. In particular, this means that $x \in A$, thus completing the proof.

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Are you clearly stating what you're assuming and what you're trying to prove?

completing the proof.

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Are you making specific claims about specific variables? Your proof should NOT have statements of the form "every element of C".

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Are all variables properly introduced and scoped? You should be able to point at every variable and say that it is either:

1) an arbitrarily chosen value - owned by the reader

2) an existentially instantiated value – owned by no one

3) an explicitly chosen value – owned by you (the proof writer)

Next Time

- First-Order Translations
 - How do we translate from English into first-order logic?
- Quantifier Orderings
 - How do you select the order of quantifiers in first-order logic formulas?
- Negating Formulas
 - How do you mechanically determine the negation of a first-order formula?
- Expressing Uniqueness
 - How do we say there's just one object of a certain type?