# First-Order Logic Part One 

## Recap from Last Time

## Recap So Far

- A propositional variable is a variable that is either true or false.
- The propositional connectives are as follows:
- Negation: $\neg p$
- Conjunction: $p \wedge q$
- Disjunction: $p \vee q$
- Implication: $p \rightarrow q$
- Biconditional: $p \leftrightarrow q$
- True: $\top$
- False: $\perp$


## Take out a sheet of paper!

## What's the truth table for the $\rightarrow$ connective?

What's the negation of $p \rightarrow q$ ?

New Stuff!

First-Order Logic

## What is First-Order Logic?

- First-order logic is a logical system for reasoning about properties of objects.
- Augments the logical connectives from propositional logic with
- predicates that describe properties of objects,
- functions that map objects to one another, and
- quantifiers that allow us to reason about multiple objects.


## Some Examples

Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka)


Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

Likes(You, Eggs) $\wedge$ Likes(You, Tomato) $\rightarrow$ Likes $($ You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

# Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka) 

Learns(You, History) v ForeverRepeats(You, History)
In(MyHeart, Havana) ^ TookBackTo(Him, Me, EastAtlanta)

These blue terms are called constant symbols. Unlike propositional variables, they refer to objects, not propositions.

Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka)
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The red things that look like function calls are called predicates.
Predicates take objects as arguments and evaluate to true or false.

Likes(You, Eggs) ^Likes(You, Tomato) $\rightarrow$ Likes(You, Shakshuka)
Learns(You, History) v ForeverRepeats(You, History)
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#### Abstract

What remains are traditional propositional connectives. Because each predicate evaluates to true or false, we can connect the truth values of predicates using normal propositional connectives.


## Reasoning about Objects

- To reason about objects, first-order logic uses predicates.
- Examples:

Cute(Quokka)
Cool(CS103 students)

- Applying a predicate to arguments produces a proposition, which is either true or false.
- Typically, when you're working in FOL, you'll have a list of predicates, what they stand for, and how many arguments they take. It'll be given separately from the formulas you write.


## First-Order Sentences

- Sentences in first-order logic can be constructed from predicates applied to objects:

$$
\begin{gathered}
\text { Cute }(a) \rightarrow \operatorname{Dikdik}(a) \vee \operatorname{Kitty}(a) \vee \operatorname{Puppy}(a) \\
\text { Succeeds }(Y o u)
\end{gathered} \leftrightarrow \operatorname{Practices(You)} \text { (Y) }
$$

$$
x<8 \rightarrow x<137
$$

The less-than sign is just another predicate. Binary predicates are sometimes written in infix notation this way.

Numbers are not "built in" to first-order logic. They're constant symbols just like
"You" and "a" above.

## Equality

- First-order logic is equipped with a special predicate $=$ that says whether two objects are equal to one another.
- Equality is a part of first-order logic, just as $\rightarrow$ and $\neg$ are.
- Examples:

$$
\begin{aligned}
& \text { MilesMorales }=\text { SpiderMan } \\
& \text { MorningStar }=\text { EveningStar }
\end{aligned}
$$

- Equality can only be applied to objects; to state that two propositions are equal, use $\leftrightarrow$.


## Let's see some more examples.

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) = StarOf(FavoriteMovieOf(Date))

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

# FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) ^ StarOf(FavoriteMovieOf(You)) = StarOf(FavoriteMovieOf(Date)) 

> FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date) $)$

# FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) ^ StarOf(FavoriteMovieOf(You)) = StarOf(FavoriteMovieOf(Date)) 

These purple terms are functions.
Functions take objects as input and produce objects as output.

> FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date) $)$

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

FavoriteMovieOf(You) $\neq$ FavoriteMovieOf(Date) $\wedge$ StarOf(FavoriteMovieOf(You)) $=$ StarOf(FavoriteMovieOf(Date))

## Functions

- First-order logic allows functions that return objects associated with other objects.
- Examples:

$$
\begin{gathered}
\text { ColorOf(Money) } \\
\text { MedianOf( } x, y, z) \\
x+y
\end{gathered}
$$

- As with predicates, functions can take in any number of arguments, but always return a single value.
- Functions evaluate to objects, not propositions.


## Objects and Predicates

- When working in first-order logic, be careful to keep objects (actual things) and propositions (true or false) separate.
- You cannot apply connectives to objects:
$\triangle \quad$ Venus $\rightarrow$ TheSun
- You cannot apply functions to propositions:
$\triangle \operatorname{StarOf(IsRed}($ Sun $) \wedge$ IsGreen(Mars)) $₫$
- Ever get confused? Just ask!


## The Type-Checking Table

|  | $\ldots$ operate on $\ldots$ | $\ldots$ and produce |
| :---: | :---: | :---: |
| Connectives <br> $(\leftrightarrow, \wedge$, etc. $) \ldots$ | propositions | a proposition |
| Predicates <br> $(=$, etc. $) \ldots$ | objects | a proposition |
| Functions $\ldots$ | objects | an object |

## One last (and major) change

## Some spider is radioactive.

## Some spider is radioactive.

## $\exists s .(S p i d e r(s) ~ \wedge ~ R a d i o a c t i v e(s)) ~$

## Some spider is radioactive.

## $\exists s .(S p i d e r(s) \wedge$ Radioactive(s))

$\exists$ is the existential quantifier and says "for some choice of s, the following is true."

## The Existential Quantifier

- A statement of the form


## $\exists x$. some-formula

is true if there exists a choice of $x$ where some-formula is true when that $x$ is plugged into it.

- Examples:
$\exists x .(E v e n(x) \wedge \operatorname{Prime}(x))$
$\exists x .($ TallerThan $(x, m e) \wedge \operatorname{LighterThan(x,~me))~}$
$(\exists w . \operatorname{Will}(w)) \rightarrow(\exists x . \operatorname{Way}(x))$


## The Existential Quantifier


$\exists x . \operatorname{Smiling}(x)$

## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



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## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



## The Existential Quantifier



$(\exists x . S m i l i n g(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier


$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier


$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier


$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat $(y))$

## The Existential Quantifier



Is this part of the statement true or false?
$(\exists x . S m i l i n g(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier



Is this part of the statement true or false?
$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat(y))

## The Existential Quantifier


$(\exists x . \operatorname{Smiling}(x)) \rightarrow(\exists y$. WearingHat $(y))$

## The Existential Quantifier


$(\exists x . S m i l i n g(x)) \rightarrow(\exists y$. WearingHat $(y))$

## Fun with Edge Cases

$\exists x . \operatorname{Smiling}(x)$

## Fun with Edge Cases

Existentially-quantified statements are false in an empty world, since nothing exists, period!
$\exists x . S m i l i n g(x)$

## Some Technical Details

## Variables and Quantifiers

- Each quantifier has two parts:
- the variable that is introduced, and
- the statement that's being quantified.
- The variable introduced is scoped just to the statement being quantified.
$(\exists x . \operatorname{Loves}(Y o u, x)) \wedge(\exists y . \operatorname{Loves}(y, Y o u))$


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The variable $x$ just lives here.

The variable $y$ just lives here.

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- the statement that's being quantified.
- The variable introduced is scoped just to the statement being quantified.
$(\exists x . \operatorname{Loves}(Y o u, x)) \wedge(\exists x . \operatorname{Loves}(x, Y o u))$

The variable $x$ just lives here.

A different variable, also named $x$, just lives here.

## Operator Precedence (Again)

- When writing out a formula in first-order logic, quantifiers have precedence just below $\neg$.
- The statement

$$
\exists x . P(x) \wedge R(x) \wedge Q(x)
$$

is parsed like this:

$$
\triangle \quad(\exists x . P(x)) \wedge(R(x) \wedge Q(x))
$$

- This is syntactically invalid because the variable $x$ is out of scope in the back half of the formula.
- To ensure that $x$ is properly quantified, explicitly put parentheses around the region you want to quantify:

$$
\exists x .(P(x) \wedge R(x) \wedge Q(x))
$$

"For any natural number $n$, $n$ is even if and only if $n^{2}$ is even"
"For any natural number $n$, $n$ is even if and only if $n^{2}$ is even"
$\forall n .\left(n \in \mathbb{N} \rightarrow\left(\operatorname{Even}(n) \leftrightarrow \operatorname{Even}\left(n^{2}\right)\right)\right)$

## "For any natural number $n$,

 $n$ is even if and only if $n^{2}$ is even"
## $\forall n .\left(n \in \mathbb{N} \rightarrow\left(\operatorname{Even}(n) \leftrightarrow \operatorname{Even}\left(n^{2}\right)\right)\right)$

$\forall$ is the universal quantifier and says "for any choice of $n$, the following is true."

## The Universal Quantifier

- A statement of the form


## $\forall x$. some-formula

is true if, for every choice of $x$, the statement some-formula is true when $x$ is plugged into it.

- Examples:
$\forall$. $($ Puppy $(p) \rightarrow$ Cute $(p))$
$\forall a$. (EatsPlants(a) v EatsAnimals(a))
Tallest(SultanKösen) $\rightarrow$
$\forall x$. (SultanKösen $\neq x \rightarrow$ ShorterThan( $x$, SultanKösen))


## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



Since Smiling(x) is true for every choice of $x$, this statement evaluates to true.
$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier



Since Smiling ( $x$ ) is true for every choice of $x$, this statement evaluates to true.
$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier


$\forall x . \operatorname{Smiling}(x)$

## The Universal Quantifier



## The Universal Quantifier



## The Universal Quantifier



Since Smiling(x) is false for this choice $x$, this statement evaluates to false.

## The Universal Quantifier



Since Smiling(x) is false for this choice $x$, this statement evaluates to false.

## Question: In this world, is the first-order logic

 statement below true or false?Respond at pollev.com/zhenglian740

$(\forall x . \operatorname{Smiling}(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier


$(\forall x . S m i l i n g(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this part of the statement true or false?
$(\forall x . S m i l i n g(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this overall statement true or false in this scenario?
$(\forall x$. Smiling $(x)) \rightarrow(\forall y$. WearingHat $(y))$

## The Universal Quantifier



Is this overall statement true or false in this scenario?

## $(\forall x . \operatorname{Smiling}(x)) \rightarrow(\forall y$. WearingHat $(y))$

## Fun with Edge Cases

$\forall x . \operatorname{Smiling}(x)$

## Fun with Edge Cases

Universally-quantified statements are said to be vacuously true in empty worlds.
$\forall x . S m i l i n g(x)$

## Let's take a quick break!

## Translating into First-Order Logic

## Translating Into Logic

- First-order logic is an excellent tool for manipulating definitions and theorems to learn more about them.
- Need to take a negation? Translate your statement into FOL, negate it, then translate it back.
- Want to prove something by contrapositive? Translate your implication into FOL, take the contrapositive, then translate it back.


## Translating Into Logic

- When translating from English into firstorder logic, we recommend that you

$$
\begin{aligned}
& \text { think of first-order logic as a } \\
& \text { mathematical programming } \\
& \text { language. }
\end{aligned}
$$

- Your goal is to learn how to combine basic concepts (quantifiers, connectives, etc.) together in ways that say what you mean.

Using the predicates

- Smiling(x), which states that $x$ is smiling, and
- WearingHat( $x$ ), which states that $x$ is wearing a hat, write a sentence in first-order logic that says


## some smiling person wears a hat.

Try it yourself: Give this your best shot - it's okay if you're not sure!

Respond at pollev.com/zhenglian740
"Some smiling person wears a hat."
$\exists x .(S m i l i n g(x) \wedge$ WearingHat(x))
$\exists x .(S m i l i n g(x) \rightarrow$ WearingHat $(x))$

"Some smiling person wears a hat."
$\exists x$. (Smiling $(x) \wedge$ WearingHat(x))
$\exists x .($ Smiling $(x) \rightarrow$ WearingHat $(x))$

"Some smiling person wears a hat."
$\exists x$. (Smiling $(x) \wedge$ WearingHat(x))
$\exists x .($ Smiling $(x) \rightarrow$ WearingHat $(x))$

"Some smiling person wears a hat." True
$\exists x$. (Smiling $(x) \wedge$ WearingHat(x))
$\exists x .($ Smiling $(x) \rightarrow$ WearingHat $(x))$

"Some smiling person wears a hat." True
$\exists x$. (Smiling $(x) \wedge$ WearingHat(x))
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"Some smiling person wears a hat." True
$\exists x$. (Smiling(x) ^ WearingHat(x)) True $\exists x .($ Smiling $(x) \rightarrow$ WearingHat $(x))$

"Some smiling person wears a hat." True
$\exists x$. (Smiling(x) ^ WearingHat(x)) True
$\exists x .($ Smiling $(x) \rightarrow$ WearingHat $(x)) \quad$ True

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$\exists x$. (Smiling $(x) \wedge$ WearingHat(x)) False
$\exists x .(S m i l i n g(x) \rightarrow$ WearingHat $(x)) \quad$ True

"Some smiling person wears a hat." False
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"Some smiling person wears a hat." False
$\exists x$. (Smiling $(x) \wedge$ WearingHat(x)) False
$\exists x .($ Smiling $(x) \rightarrow$ WearingHat $(x))$ True

# "Some $P$ is a $Q "$ 

translates as

## ヨx. (P(x) ^ Q(x))

## Useful Intuition:

Existentially-quantified statements are false unless there's a positive example.

## ヨx. (P(x) ^ Q(x))

If $x$ is an example, it must have property $P$ on top of property $Q$.

Using the predicates

- Smiling(x), which states that $x$ is smiling, and
- WearingHat( $x$ ), which states that $x$ is wearing a hat, write a sentence in first-order logic that says


## every smiling person wears a hat.

Try it yourself: Give this your best shot - it's okay if you're not sure!

Respond at pollev.com/zhenglian740
"Every smiling person wears a hat."
$\forall x$. (Smiling $(x) \wedge$ WearingHat $(x))$
$\forall x .($ Smiling $(x) \rightarrow$ WearingHat $(x))$

"Every smiling person wears a hat."
$\forall x .(S m i l i n g(x) \wedge$ WearingHat(x)) $\forall x .(S m i l i n g(x) \rightarrow$ WearingHat $(x))$

"Every smiling person wears a hat." True $\forall x$. (Smiling $(x) \wedge$ WearingHat(x)) $\forall x .(S m i l i n g(x) \rightarrow$ WearingHat $(x))$

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"Every smiling person wears a hat." True $\forall x$. (Smiling $(x)$ ^ WearingHat $(x)$ ) False $\forall x .($ Smiling $(x) \rightarrow$ WearingHat $(x)) \quad$ True

## "All P's are Q's"

translates as
$\forall x .(P(x) \rightarrow Q(x))$

## Useful Intuition:

Universally-quantified statements are true unless there's a counterexample.

## $\forall x .(P(x) \rightarrow Q(x))$

If $x$ is a counterexample, it must have property $P$ but not have property $Q$.

## Good Pairings

- The $\forall$ quantifier usually is paired with $\rightarrow$.

$$
\forall x .(P(x) \rightarrow Q(x))
$$

- The $\exists$ quantifier usually is paired with $\wedge$.

$$
\exists x .(P(x) \wedge Q(x))
$$

- In the case of $\forall$, the $\rightarrow$ connective prevents the statement from being false when speaking about some object you don't care about.
- In the case of $\exists$, the $\wedge$ connective prevents the statement from being true when speaking about some object you don't care about.

Proofwriting Workshop

## An Incorrect Set Theory Proof

Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
$\triangle$ Incorrect! $\triangle$ Proof: Consider arbitrary sets $A$, $B$, and $C$ where $C \subseteq A \cup B$.

This means that every element of $C$ is in either $A$ or $B$. If all elements of $C$ are in $A$, then $C \subseteq A$. Alternately, if everything in $C$ is in $B$, then $C \subseteq B$. In either case, everything inside of $C$ has to be contained in at least one of these sets, so the theorem is true.

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This is just repeating definitions and not making specific claims about specific variables.

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Why is this bad?

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While this claim is true, it does not imply the theorem is true. In fact, this theorem is actually false.

## Let's Draw Some Pictures!

Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.

## Let's Draw Some Pictures!

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## Let's Draw Some Pictures!

Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.

Recall the intuition of a subset being "something I can circle"


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So $\mathbf{C} \subseteq A$ would mean that $C$ is something I can circle in this region.

A


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Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.

Recall the intuition of a subset being "something I can circle"

Likewise, $\mathbf{C} \subseteq \mathrm{B}$ would mean that C is something I can circle in this region.

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But when I look at $A \cup B$, I can draw $C$ as a circle containing elements from both $A$ and $B$ !


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But when I look at $A \cup B$, I can draw $C$ as a circle containing elements from both $A$ and $B$ !


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Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.

Using this visual intuition, come up with a choice of sets $A, B$, and $C$ that show this claim is false.
Respond at pollev.com/zhenglian740


## Proofs vs. Disproofs

- A proof is an argument that explains why some theorem is true.
- A disproof is an argument that explains why some claim is false.
- You've seen lots of examples of proofs. What does a disproof look like?

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Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \& A$ and $C \& B$.

Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \not \subset A$ and $C \not \subset B$. Consider the sets $A=\{1\}$


Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \not \subset A$ and $C \not \subset B$. Consider the sets $A=\{1\}, B=\{2\}$


Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \notin A$ and $C \notin B$. Consider the sets $A=\{1\}, B=\{2\}$, and $C=\{1,2\}$.


Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \not \subset A$ and $C \not \subset B$. Consider the sets $A=\{1\}, B=\{2\}$, and $C=\{1,2\}$. Now notice that $\{1,2\} \subseteq A \cup B$ so $C \subseteq A \cup B$


Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \not \subset A$ and $C \nsubseteq B$. Consider the sets $A=\{1\}, B=\{2\}$, and $C=\{1,2\}$. Now notice that $\{1,2\} \subseteq A \cup B$ so $C \subseteq A \cup B$, but $C \& A$ because $2 \in C$ but $2 \notin A$


Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \notin A$ and $C \notin B$. Consider the sets $A=\{1\}, B=\{2\}$, and $C=\{1,2\}$. Now notice that $\{1,2\} \subseteq A \cup B$ so $C \subseteq A \cup B$, but $C \notin A$ because $2 \in C$ but $2 \notin A$, and $C \& B$ because $1 \in C$ but $1 \notin B$.


Claim: If $A, B$, and $C$ are sets and $C \subseteq A \cup B$, then $C \subseteq A$ or $C \subseteq B$.
Disproof: We will show that there are sets $A, B$, and $C$ where $C \subseteq A \cup B$, but $C \notin A$ and $C \notin B$. Consider the sets $A=\{1\}, B=\{2\}$, and $C=\{1,2\}$. Now notice that $\{1,2\} \subseteq A \cup B$ so $C \subseteq A \cup B$, but $C \nsubseteq A$ because $2 \in C$ but $2 \notin A$, and $C \& B$ because $1 \in C$ but $1 \notin B$.
Thus we've found a set $C$ which is a subset of $A \cup B$ but is not a subset of either $A$ or $B$, which is what we needed to show. ■

## Proofwriting Advice

- Be very wary of proofs that speak generally about "all objects" of a particular type.
- As you've just seen, it's easy to accidentally prove a false statement at this level of detail.
- Making broad, high-level claims often indicates deeper logic errors or conceptual misunderstanding (like code smell but for proofs!)


## Proofwriting Advice

A Very Good Idea: After you've written a draft of a proof, run through all of the points on the Proofwriting Checklist.

- This is a great exercise that you can do with a partner!


## Proofs on Subsets

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

## Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then

 $C \subseteq A$ and $C \subseteq B$.Hold on, isn't this the claim we just disproved?

## Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then

 $C \subseteq A$ and $C \subseteq B$.Notice that that's an intersection, not a union! It turns out that this claim is actually true.

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Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Recall the intuition of a subset being "something I can circle"

When I look at $\mathrm{A} \cap \mathrm{B}$, any circle I can draw in this region can be found in both $A$ and $B$.


## Let's Draw Some Pictures!

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.


# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

What We're Assuming
What We Need To Show

When confronted with a theorem to prove, the first step is to make sure you
understand where you're starting and where you're going.

# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

What We're Assuming

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- A, B, and C are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$

A great proofwriting strategy is to write down relevant definitions.
This gives you a better sense of what you need to prove and what tools you have at hand.

# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

What We're Assuming

- $A, B$, and $C$ are sets

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

Before we start:

- What is the definition of subset?
- How do you prove that one set is a subset of another?
- If you know that one set is a subset of another, what can you conclude?

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- $A, B$, and $C$ are sets

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- $C \subseteq A \cap B$

Definition: If $S$ and $T$ are sets, then $S \subseteq T$ when for every $x \in S$, we have $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- $A, B$, and $C$ are sets
- $C \subseteq A \cap B$

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$

Definition: If $S$ and $T$ are sets, then $S \subseteq T$ when for every $x \in S$, we have $x \in T$.

To prove that $S \subseteq T$ :
Pick an arbitrary $x \in S$, then prove $x \in T$.
If you know that $S \subseteq T$ :
If you have an $x \in S$, you can conclude $x \in T$.

# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

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## Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

This reading of the definition is usually helpful for unpacking this

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$

Definition: If $S$ and $T$ are sets, then $S \subseteq T$ when for every $x \in S$, we have $x \in T$.

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# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

What We're Assuming

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$
- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$


# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

What We're Assuming

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

> Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$

How can we apply this general template to our specific problem?

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- A, B, and C are sets
- $C \subseteq A \cap B$

> Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- Pick an $x \in C$, show that $x$ $\in A$
- Pick an $x \in C$, show that $x$ $\in B$

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- A, B, and C are sets
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> Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$


## What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- Pick an $x \in C$, show that $x$ $\in A$
- Pick an $x \in C$, show that $x$ $\in B$

Now we know that ultimately, we're going to have to do these two things. Let's see what tools we have that can get us here!

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- A, B, and C are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

> Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$


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- $C \subseteq A$ and $C \subseteq B$
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What We're Assuming

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This reading of the definition is usually helpful for unpacking this column!


Definition: If $S$ and $T$ ard sets, then $S \subseteq T$ when for every $x \in S$, whave $x \in T$.

Pick an arbitrary $x \in S$, th n prove $x \in T$.
If you know that $S \subseteq T$ :
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## Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then

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- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

> Our Tools

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$


## What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- Pick an $x \in C$, show that $x$ $\in A$
- Pick an $x \in C$, show that $x$ $\in B$

Before we continue:

- What is the definition of set intersection?


# Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$. 

What We're Assuming

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$


## What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- Pick an $x \in C$, show that $x$ $\in A$

Definition: The set $S \cap T$ is the set where, for any $\chi$ : $x \in S \cap T \quad$ when $\quad x \in S$ and $x \in T$

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

## What We Need To Show

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$

Definition: The set $S \cap T$ is the set where, for any $x$ :

$$
x \in S \cap T \quad \text { when } \quad x \in S \text { and } x \in T
$$

To prove that $x \in S \cap T$ :
Prove both that $x \in S$ and that $x \in T$.
If you know that $x \in S \cap T$ : You can conclude both that $x \in S$ and that $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- A, B, and C are sets
- $C \subseteq A \cap B$

```
Definition: The set S \capT is the set where, for any }x\mathrm{ : \(x \in S \cap T \quad\) when \(\quad x \in S\) and \(x \in T\)
To prove that \(x \in S \cap T\) : Prove both that \(x \in S\) ad that \(x \in T\).
If you know that \(x \in S \cap T\) : You can conclude both that \(x \in S\) and that \(x \in T\).
```

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$
Our Tools
- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$.
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- Pick an $x \in C$, show that $x$ $\in A$
- Pick an $x \in C$, show that $x$ $\in B$

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

What We're Assuming

- $A, B$, and $C$ are sets
- $\mathrm{C} \subseteq \mathrm{A} \cap \mathrm{B}$
Our Tools
- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$.
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

What We Need To Show

- $C \subseteq A$ and $C \subseteq B$
- Pick an $x \in C$, show that $x$ $\in A$
- Pick an $x \in C$, show that $x$ $\in B$

Let's go and try and do the proof with what we've got here!

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$


## Relevant Definitions

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$
- Conclude $x \in A$


## Relevant Definitions

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$



## Relevant Definitions

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$
- $x \in A \cap B$
- Conclude $x \in A$

Relevant Definitions

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
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- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
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- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving C $\subseteq A$
- Pick an $x \in C$
- $x \in A \cap B$
- $x \in A$ and $x \in B$
- Conclude $x \in A$

Relevant Definitions

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
- If you know that $x \in S \cap T$, we can conclude that $x \in S$ and $x \in T$.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$
- $x \in A \cap B$
- $x \in A$ and $x \in B$
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## Relevant Definitions

- In general to show that $S \subseteq T$, pick an arbitrary $x \in S$, show that $x \in T$
- If you know that $S \subseteq T$ and you have an $x \in S$, you can conclude $x \in T$.
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Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$
- $x \in A \cap B$
- $x \in A$ and $x \in B$
- Conclude $x \in A$

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$
- $x \in A \cap B$
- $x \in A$ and $x \in B$
- Conclude $x \in A$

We also need to prove that $C \subseteq B$.

Notice that if you take the outline here and literally swap the variable A for the variable $B$, you get a working proof.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq B \cap A$
- Proving $C \subseteq B$
- Pick an $x \in C$
- $x \in B \cap A$
- $x \in B$ and $x \in A$
- Conclude $x \in B$

In a case like this where your proof would have two completely symmetric branches, it's fine to write up just one and say
"by symmetry, [the other branch]
is also true."

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Rough Outline

- Assume $C \subseteq A \cap B$
- Proving $C \subseteq A$
- Pick an $x \in C$
- $x \in A \cap B$
- $x \in A$ and $x \in B$
- Conclude $x \in A$

Try it yourself: Take a few minutes and write up a proof of the theorem using this outline.

Then share your proof with your neighbors and critique each other!

Respond at pollev.com/zhenglian740

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Proof: Let $A, B$, and $C$ be arbitrary sets where $C \subseteq A \cap B$. We need to show that $C \subseteq A$ and $C \subseteq B$. Because the roles of $A$ and $B$ in this proof are symmetric, we can just prove that $C \subseteq A$.
Choose any element $x \in C$. Since $C \subseteq A \cap B$, we know that $x \in A \cap B$. This tells us that $x \in A$ and $x \in B$. In particular, this means that $x \in A$, thus completing the proof.

Theorem: If $A, B$, and $C$ are sets and $C \subseteq A \cap B$, then $C \subseteq A$ and $C \subseteq B$.

Proof: Let $A, B$, and $C$ be arbitrary sets where $C \subseteq A \cap B$. We need to show that $C \subseteq A$ and $C \subseteq B$. Because the roles of $A$ and $B$ in this proof are

Are you clearly stating what you're assuming and what you're trying to prove?
completing the proof.

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Proof: Let $A, B$, and $C$ be arbitrary sets where $C \subseteq A \cap B$. We need to show that $C \subseteq A$ and $C \subseteq B$. Because the roles of $A$ and $B$ in this proof are symmetric, we can just prove that $C \subseteq A$.
Choose any element $x \in C$. Since $C \subseteq A \cap B$, we know that $x \in A \cap B$. This tells us that $x \in A$ and $x \in B$. Tn narticular this means that $x \in A$ thus

Are you making specific claims about specific variables? Your proof should NOT have statements of the form "every element of $C$ ".

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Are all variables properly introduced and scoped? You should be able to point at every variable and say that it is either:

1) an arbitrarily chosen value - owned by the reader
2) an existentially instantiated value - owned by no one
3) an explicitly chosen value - owned by you (the proof writer)

## Next Time

- First-Order Translations
- How do we translate from English into first-order logic?
- Quantifier Orderings
- How do you select the order of quantifiers in first-order logic formulas?
- Negating Formulas
- How do you mechanically determine the negation of a first-order formula?
- Expressing Uniqueness
- How do we say there's just one object of a certain type?

